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On bandwidth, cutwidth, and quotient graphs


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ON BANDWIDTH, CUTWIDTH, AND QUOTIENT GRAPHS (*)

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Abstract. – The bandwidth and the cutwidth are fundamental parameters related to many problems modeled in terms of graphs. In this paper, we present a general method for finding upper bounds of the bandwidth of a graph from the ones the quotient graph and induced subgraphs issued from any of its partitions, as well as an equivalent result for cutwidth. Moreover, general lower bounds are obtained by using vertex- and edge-bisection notions.

These results are applied, in a second time, to several interconnection networks. By choosing convenient vertex partitions and judicious internal numberings of the vertices of the partition blocks, we prove in this paper original bounds for the binary de Bruijn and Butterfly graphs, and summarize results for the Shuffle-Exchange, FFT, and CCC graphs.

Résumé. – Les largeurs de bande et de coupe sont deux paramètres fondamentaux qui interviennent dans la formulation de nombreux problèmes modélisés en termes de graphes.

Dans cet article, nous présentons une méthode permettant de majorer la largeur de bande d’un graphe en fonction de celle du graphe quotient et des sous-graphes induits par n’importe laquelle de ses partitions, ainsi qu’un résultat équivalent pour la largeur de coupe. De plus, nous proposons des minorations générales de ces paramètres à partir des bissections sommet et arête du graphe.

Ces résultats sont appliqués à plusieurs réseaux d’interconnexion. En choisissant des partitions adaptées et des numérotations judicieuses des sommets des parties de ces partitions, nous prouvons des majorations originales pour les graphes de de Bruijn binaire et Butterfly, et récapitulons nos résultats pour les graphes Shuffle-Exchange, FFT, et CCC.

1. INTRODUCTION

In all the paper, $G$ is a graph with $n$ vertices, the vertex and edge sets of which are denoted $V(G)$ and $E(G)$ respectively. An ordering $\varphi$ of $G$ is a one-to-one mapping of $V(G)$ in the set $\{0, 1, \ldots, (n - 1)\}$, and $\Phi(G)$ is the set of all orderings $\varphi$ of $G$. 

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The ordering of graphs is equivalent to their embedding into the path graph of same order, which allows us to define the dilation \( \text{dil}(\varphi) \) and congestion \( \text{cg}(\varphi) \) of an ordering \( \varphi \) as the ones of the corresponding embedding:

\[
\text{dil}(\varphi) \overset{\text{def}}{=} \max_{\{v', v''\} \in E(G)} (|\varphi(v') - \varphi(v'')|), \\
\text{cg}(\varphi) \overset{\text{def}}{=} \max_{v \in V(G)} (|\{v', v''\} \in E(G)/\varphi(v') \leq \varphi(v) < \varphi(v'')|).
\]

The bandwidth of \( G \) is then defined as

\[
\text{Bd}(G) \overset{\text{def}}{=} \min_{\varphi \in \Phi(G)} (\text{dil}(\varphi)),
\]

and the cutwidth of \( G \) is

\[
\text{Cw}(G) \overset{\text{def}}{=} \min_{\varphi \in \Phi(G)} (\text{cg}(\varphi)).
\]

The bandwidth and the cutwidth are two formal parameters of graphs which appear in the formulation of many problems modeled in terms of graphs. For instance, bandwidth has been used in code correction [14], gaussian elimination for sparse matrices [6, 12], and V.L.S.I. layout [18, 30], and cutwidth in V.L.S.I. design [9, 20, 25, 29].

General bounds have been proven for bandwidth in [6, 7], but since it has been shown in [10, 11, 12] that the problem of finding the bandwidth and the cutwidth of any given graph is NP-complete, the main way to compute them is by the means of heuristics, as in [12, 13, 28].

In several occasions [7, 12], the method used to compute bandwidth involves a two-stage process, where vertices are clustered before being numbered within these clusters. Formalizing this approach by the means of quotient graphs, we present in this paper a method for finding upper bounds of the bandwidth and cutwidth of a graph, based on the bandwidth and cutwidth of its quotient graph and induced subgraphs with respect to the blocks of partitions of \( V(G) \). All the necessary definitions and notations are given in Section 2.

This method leads to general upper bounds, stated in Theorems 1 and 5. When the structures of the graphs are convenient, the vertex orderings induced by the partitions can be finely tuned to improve the upper bounds; in that sense, our results can be seen as a framework for finding good
orderings rather than plain theorems. Lower bounds are also proven in Propositions 3 and 7, using the vertex- and edge-bisections of \( G \). Section 3 contains all these results.

In Section 4, we use these results to prove original bounds for the binary de Bruijn and Butterfly graphs. Equivalent results for the Shuffle-Exchange, FFT, and CCC graphs are also summarized.

2. DEFINITIONS AND NOTATIONS

It is assumed that the reader is familiar with standard graph theoretic notations; see [5] for reference. In particular, \( \text{diam}(G) \) is the diameter of graph \( G \), \( \delta(v) \) denotes the degree of a vertex \( v \), and \( \delta(G) \overset{\text{def}}{=} \min_{v \in V(G)} \delta(v) \) and \( \Delta(G) \overset{\text{def}}{=} \max_{v \in V(G)} \delta(v) \) are the minimum and maximum degrees of graph \( G \), respectively.

In the following, \( a \mod b \) denotes the remainder of the euclidian division of \( a \) by \( b \), and \( \oplus \) the exclusive-or operator on binary representations of integer numbers.

2.1. Partitions

Given a non-empty \( n \)-vertex set \( V(G) \) of a graph \( G \), a partition \( \Pi \) of \( V(G) \) is a family of \( N \) non-empty mutually disjoint subsets called blocks, the union of which is equal to \( V(G) \). We denote \( \mathcal{P}(G) \) the set of all the partitions of \( V(G) \).

For all vertices \( v \) in \( V(G) \), \( \pi(v) \) denotes the block of \( \Pi \) containing \( v \). The size of the biggest block of a partition is \( \max_{\pi} \overset{\text{def}}{=} \max_{\pi \in \Pi} (|\pi|) \).

For all blocks \( \pi \) in \( \Pi \), \( \omega(\pi) \) is the cocycle of \( \pi \) in \( G \), i.e. the edge set \( \{v', v''\} \in E(G) / v' \in \pi, v'' \notin \pi \}. \) The size of the biggest cocycle in a partition is \( \max_{\omega} \overset{\text{def}}{=} \max_{\pi \in \Pi} (|\omega(\pi)|) \). \( E_I(\pi) \) and \( E_E(\pi) \) respectively denote the sets of edges which have both ends in \( \pi \) (i.e. internal edges) and none of their ends in \( \pi \) (i.e. external edges). These notations are extended to the whole partition, with \( E_I(\Pi) \overset{\text{def}}{=} \bigcup_{\pi \in \Pi} (E_I(\pi)) \) and \( E_E(\Pi) \overset{\text{def}}{=} E(G) - E_I(\Pi) \).

2.2. Quotient graphs

For all blocks \( \pi \) in \( \Pi \), \( G[\pi] \) denotes the subgraph of \( G \) induced by \( \pi \). The quotient graph \( Q = G/\Pi \) of a graph \( G \) with respect to a partition \( \Pi \) of \( V(G) \)
is the graph such that $V(Q) = \Pi$, and \{\pi', \pi''\} belongs to $E(Q)$ if and only if there exist $v'$ in $\pi'$ and $v''$ in $\pi''$ such that \{$v', v''$\} belongs to $E(G)$.

2.3. Orderings

An ordering $\varphi$ of a graph $G$ is a one-to-one mapping of $V(G)$ in \{0, 1, ..., $(n - 1)$\}. The restriction of $\varphi$ to a block $\pi$ of $G$, denoted $\varphi|_{\pi}$, is the ordering of $G[\pi]$ such that for all vertices $v'$ and $v''$ in $\pi$, $(\varphi(v') < \varphi(v'')) \Leftrightarrow (\varphi|_{\pi}(v') < \varphi|_{\pi}(v''))$.

$\varphi$ defines an implicit orientation of the edges of $G$, which allows us to define the lower and upper half-degrees of any vertex:

$$\delta^-_\varphi(v) \text{ def } \{|\{v', v\} \in E(G)/\varphi(v') < \varphi(v)\|,$$

and

$$\delta^+_\varphi(v) \text{ def } \{|\{v, v''\} \in E(G)/\varphi(v) < \varphi(v'')\|.$$

Let $\Phi_Q(G, \Pi)$ be the set of all orderings $\varphi_Q$ of $Q = G/\Pi$. For all blocks $\pi$ of $\Pi$ ordered by $\varphi_Q$, we define the numbers

$$\delta^-_{\Pi, \varphi_Q}(\pi) \text{ def } \min_{v \in \pi} (\{|\{v', v\} \in \omega(\pi(v))/\varphi_Q(\pi(v')) < \varphi_Q(\pi(v))\|),$$

and

$$\delta^+_{\Pi, \varphi_Q}(\pi) \text{ def } \min_{v \in \pi} (\{|\{v, v''\} \in \omega(\pi(v))/\varphi_Q(\pi(v'')) > \varphi_Q(\pi(v))\|).$$

They are the minimum number of edges which link any vertex of block $\pi$ to vertices belonging to blocks of smaller and bigger numbers, respectively.

Let $G$ be a graph, $\Pi$ a partition of $V(G)$ into $N$ blocks, and $\varphi_Q$ an ordering in $\Phi_Q(G, \Pi)$. An ordering $\varphi_G$ of $V(G)$ is compatible with $\Pi$ and $\varphi_Q$ if and only if vertices in blocks of increasing numbers have increasing numbers, i.e.:

$$\forall \pi', \pi'' \in \Pi,$$

$$((\varphi_Q(\pi')) < \varphi_Q(\pi'')) \Leftrightarrow (\forall v' \in \pi', \forall v'' \in \pi'', \varphi_G(v') < \varphi_G(v'')).$$
which amounts to the following set of properties:

\[
\begin{align*}
\bigcup_{v \in \pi} \{\varphi(G)(v)\} &= (p_{\varphi_{Q}}(0), \ldots, P_{\varphi_{Q}}(\pi)) & \text{the numbers in each block form an interval;} \\
\bigcup_{t=0}^{N-1} \{p_t, \ldots, P_t\} &= \{0, \ldots, n-1\} & \text{the intervals are disjoint and cover \{0, \ldots, n-1\};} \\
\forall i \in \{1, \ldots, N-1\}, p_i &= P_{i-1} + 1 & \text{partitions of increasing numbers give intervals of increasing extrema.}
\end{align*}
\]

In particular, \(p_0 = 0, P_0 = |\varphi_{Q}^{-1}(0)| - 1, p_1 = |\varphi_{Q}^{-1}(0)|, P_1 = |\varphi_{Q}^{-1}(0)| + |\varphi_{Q}^{-1}(1)| - 1, \ldots, P_{N-1} = n - 1\). We denote \(\Phi(G, \Pi, \varphi_Q)\) the set of all these orderings. We call configuration of \(G\) any triplet \((\Pi, \varphi_Q, \varphi_G)\) such that \(\Pi\) belongs to \(\mathcal{P}(G)\), \(\varphi_Q\) belongs to \(\Phi_Q(G, \Pi)\), and \(\varphi_G\) belongs to \(\Phi_G(G, \Pi, \varphi_Q)\).

Since the partition \(\Pi \approx V(G)\) with \(n\) blocks always allows orderings \(\varphi_G \approx \varphi_Q\) which achieve the bandwidth and the cutwidth of \(G\), we have

\[
\begin{align*}
Bd(G) &= \min_{(\Pi, \varphi_Q, \varphi_G)} (\text{dil}(\varphi_G)), \\
Cw(G) &= \min_{(\Pi, \varphi_Q, \varphi_G)} (\text{cg}(\varphi_G)).
\end{align*}
\]

2.4. Bisections

The edge-bisection \(\text{bis}_e(G)\) of a graph \(G\) is the size of the edge-set of minimum cardinality whose removal splits \(G\) into two subgraphs the vertex-cardinalities of which differ by at most one.

The vertex-bisection \(\text{bis}_v(G)\) of a graph \(G\) is the size of the vertex-set of minimum cardinality whose removal splits \(G\) into two subgraphs of same vertex-cardinality. In some parts of our computations, we will be more
interested in the minimality of the size of a disconnecting vertex-set rather than in the strict equality between the sizes of the two resulting subgraphs. Therefore, we define the almost-vertex-bisection $\operatorname{bis}_v'(G)$ of a graph $G$ to be the size of the vertex-set of minimum cardinality whose removal splits $G$ into two subgraphs the vertex-cardinalities of which differ by at most one. Therefore,

$$\operatorname{bis}_v(G) = \begin{cases} 
\operatorname{bis}_v'(G) & \text{if } n - \operatorname{bis}_v'(G) \text{ is even} \\
\operatorname{bis}_v'(G) + 1 & \text{if } n - \operatorname{bis}_v'(G) \text{ is odd.}
\end{cases}$$

3. GENERAL RESULTS

In this section, we prove upper and lower bound theorems, both for bandwidth and cutwidth. The proofs of the upper bounds are based on majorations performed when considering the ordering of the vertices of the original graph with respect to the one of any of its quotient graphs. The proofs of the lower bounds are based on the almost-vertex- and edge-bisections, respectively.

**Theorem 1:** Let $G$ be a graph, $Q$ the quotient graph obtained from a partition $\Pi$ of $V(G)$, $\varphi_Q$ an ordering of $Q$ achieving the bandwidth of $Q$. For all $\varphi_G \in \Phi_G(G, \Pi, \varphi_Q)$, we have

$$\operatorname{Bd}(G) \leq \max(\varepsilon_i(\varphi_G), (\operatorname{Bd}(Q) - 1) \max_{\pi} + \varepsilon_e(\varphi_G))$$

where

$$\varepsilon_i(\varphi_G) = \max_{\pi \in \Pi} (\operatorname{dil}(\varphi_G|\pi)),
\varepsilon_e(\varphi_G) = \max_{\{v', v''\} \in E_G(\Pi)} ((P_{\varphi_G}(\pi(v')) - \varphi_G(v'))
\quad \varphi_G(v') < \varphi_G(v'') + (\varphi_G(v'') - p_{\varphi_Q}(\pi(v'')) + 1)).$$

**Proof:** We know that, for all $\varphi_G$ in $\Phi_G(G, \Pi, \varphi_Q)$, $\operatorname{Bd}(G) \leq \operatorname{dil}(\varphi_G)$. In order to make parameters of $Q$ appear within the expression of $\operatorname{dil}(\varphi_G)$, we rewrite this latter with respect to the blocks of the partition, splitting it into two terms accounting respectively for the edges of $G$ internal to the blocks and for the other edges.
For all configurations \((\Pi, \varphi_Q, \varphi_G)\), \(\text{dil}(\varphi_G)\) rewrites into:

\[
\text{dil}(\varphi_G) = \max \left( \max_{\{v', v''\} \in E_E(\Pi)} (|\varphi_G(v') - \varphi_G(v'')|), \right.
\]

\[
\left. \max_{\{v', v''\} \in E_E(\Pi)} (|\varphi_G(v') - \varphi_G(v'')|) \right).
\]

Let us evaluate these two terms separately.

- Since \(\varphi_G\) is an ordering compatible with \(\varphi_Q\), the vertices in each block are numbered in sequence, so

\[
\max_{\{v', v''\} \in E_E(\Pi)} (|\varphi_G(v') - \varphi_G(v'')|)
\]

\[
= \max_{\pi \in \Pi} \left( \max_{\{v', v''\} \in E_E(\pi)} (|\varphi_G|_{\pi}(v') - \varphi_G|_{\pi}(v'')) \right)
\]

\[
= \max_{\pi \in \Pi} (\text{dil}(\varphi_G|_{\pi})) = \varepsilon_e(\varphi_G).
\]

- When \(E_E(\Pi)\) is not empty, \(\text{Bd}(Q) \geq 1\), and for all \(\{v', v''\} \in E_E(\Pi)\), with \(q' = \varphi_Q(\pi(v')) < \varphi_Q(\pi(v'')) = q''\), we have \(q'' - q' \leq \text{Bd}(Q)\) and

\[
\varphi_G(v'') - \varphi_G(v')
\]

\[
= (\varphi_G(v'') - p_{q''} + 1) + (P_{q''-1} - p_{q'+1} + 1) + (P_{q'} - \varphi_G(v'))
\]

\[
= \sum_{q < q < q''} |\varphi_Q^{-1}(q)| + (P_{q'} - \varphi_G(v')) + (\varphi_G(v'') - p_{q''}) + 1
\]

\[
\leq (\text{Bd}(Q) - 1) \max_{\pi} + (P_{q'} - \varphi_G(v')) + (\varphi_G(v'') - p_{q''}) + 1.
\]

Therefore,

\[
\max_{\{v', v''\} \in E_E(\Pi)} (|\varphi_G(v') - \varphi_G(v'')|) \leq (\text{Bd}(Q) - 1) \max_{\pi} + \varepsilon_e(\varphi_G).
\]

Combining these two results leads to the claimed result. \(\square\)

Although the majorations performed within the proof seem rough, this theorem usually gives good results as most partitions are taken with blocks having the same size. If it is not the case, a more specific study has to be carried on.

From the above theorem, we can deduce the following straightforward corollary.
Corollary 2: If $G$ is a graph and $Q$ is the quotient graph obtained from a partition $\Pi$ of $V(G)$, then

$$\text{Bd}(G) \leq (\text{Bd}(Q) + 1) \max - 1.$$  

Proof: Notice that $\varepsilon_i(\varphi_G) \leq \max - 1$ and $\varepsilon_e(\varphi_G) \leq 2\max - 1$. □

Proposition 3: For all graphs $G$,

$$\text{Bd}(G) \geq \text{bis}'(G).$$  

Proof: Let $\varphi_G$ be an ordering of $G$ which achieves the bandwidth of $G$, i.e. such that $\max \{\varphi_G(v') - \varphi_G(v'')\} = \text{Bd}(G)$. Let us define

$$b_1 = \left\lfloor \frac{|V(G)| - \text{Bd}(G)}{2} \right\rfloor \quad \text{and} \quad b_2 = \left\lceil \frac{|V(G)| + \text{Bd}(G)}{2} \right\rceil,$$

and let $V_d = \{v \in V(G) / b_1 \leq \varphi_G(v) < b_2\}$. By what precedes, there is no edge $\{v', v''\} \in E(G)$ such that $\varphi_G(v') < b_1$ and $\varphi_G(v'') \geq b_2$, and $V_d$ disconnects $G$ into two subgraphs $G_1$ and $G_2$ such that $|V(G_1)| = |V(G_2)|$ or $|V(G_1)| = |V(G_2)| - 1$, so $\text{bis}'(G) \leq |V_d| = \text{Bd}(G)$.

Since the vertex-bisection is not known for many graphs, the most direct way to compute lower bounds for bandwidth is by means of Chung’s theorem.

Theorem 4: Chung [7]

$$\text{Bd}(G) \geq \left\lfloor \frac{n - 1}{\text{diam}(G)} \right\rfloor.$$  

The most interesting point of this bound is that it does not require a deep knowledge of the studied graph, while the (almost-) vertex-bisection is often hard to compute. As a matter of fact, the lower bounds of the bandwidth of the two chosen examples are computed using Chung’s theorem.

Let us now deal with the cutwidth.

Theorem 5: If $G$ is a graph, $Q$ the quotient graph obtained from a partition $\Pi$ of $V(G)$, and $\varphi_Q$ an ordering of $Q$ achieving the cutwidth of $Q$, then

$$\text{Cw}(G) \leq \left(\text{Cw}(Q) - \left\lfloor \frac{\delta(Q)}{2} \right\rfloor \right) \cdot \max_{\omega}$$

$$+ \max_{\pi \in \Pi} (\text{Cw}(G[\pi]) + |\omega(\pi)| - |\pi| \cdot \min (\delta^-_{\Pi, \varphi_Q}(\pi), \delta^+_{\Pi, \varphi_Q}(\pi))).$$
Proof: We know that, for all $\varphi_G$ in $\Phi_G (G, \Pi, \varphi_Q)$, $Cw (G) \leq cg (\varphi_G)$. In order to make parameters of $Q$ appear within the expression of $cg (\varphi_G)$, we rewrite this latter with respect to the blocks of the given partition, splitting it into three terms accounting respectively for the edges of $G$ external to a given block, the edges which have exactly one end in this block, and the edges internal to the block, as shown in Figure 1.

\begin{figure}
\centering
\includegraphics[width=0.5\textwidth]{figure1}
\caption{External (dashed), cocycle (dotted), and internal (solid) edges with respect to block \( \pi \).}
\end{figure}

Let

\[
\begin{align*}
f_E (\pi) &= \max_{v \in \pi} (|\{\{v', v''\} \in E_G (\pi) / \varphi_G (v') \leq \varphi_G (v) < \varphi_G (v''\})|), \\
f_C (\pi) &= \max_{v \in \pi} (|\{\{v', v''\} \in \omega (\pi) / \varphi_G (v') \leq \varphi_G (v) < \varphi_G (v''\})|),
\end{align*}
\]

and

\[
\begin{align*}
f_I (\pi) &= \max_{v \in \pi} (|\{\{v', v''\} \in E_I (\pi) / \varphi_G (v') \leq \varphi_G (v) < \varphi_G (v''\})|),
\end{align*}
\]

so that $cg (\varphi_G) \leq \max_{\pi \in \Pi} (f_E (\pi) + f_C (\pi) + f_I (\pi))$. Let us study these three terms separately. Let $\pi$ in $\Pi$ and $v$ in $\pi$.

Since $f_E (\pi)$ accounts for the edges external to $\pi$, the choice of $v$ in $\pi$ is not significant. Let us then consider the quotient graph $Q = G/\Pi$. Let $\pi^-$ and $\pi^+$ be the vertices of $Q$ such that $\varphi_Q (\pi^-) = \varphi_Q (\pi) - 1$ and $\varphi_Q (\pi^+) = \varphi_Q (\pi) + 1$, if they exist. The number of edges of $Q$ whose end numbers strictly bracket $\varphi_Q (\pi)$, which we denote $g_E (\pi)$, can be written in two different ways:

\[
g_E (\pi) = |\{\{w', w''\} \in E (Q) / \varphi_Q (w') \leq \varphi_Q (\pi^-) < \varphi_Q (w'')\}| - \delta_{\varphi_Q} (\pi),
\]
and

\[ g_E(\pi) = |\{\{\omega', \omega''\} \in E(Q)/\varphi_Q(\omega') \leq \varphi_Q(\omega'')\}| - \delta_{\varphi_Q}^+(\pi). \]

Since the first term of each right member of the above equations is, by definition, bounded by \( cg(\varphi_Q) \), we have

\[ g_E(\pi) \leq cg(\varphi_Q) - \min(\delta_{\varphi_Q}^-(\pi), \delta_{\varphi_Q}^+(\pi)), \]

and since \( \delta_{\varphi_Q}^-(\pi) + \delta_{\varphi_Q}^+(\pi) = \delta(\pi) \), the minimum of the partial degrees is always greater than or equal to \( \left[ \frac{\delta(\pi)}{2} \right] \), so \( g_E(\pi) \leq cg(\varphi_Q) - \left[ \frac{\delta(\pi)}{2} \right] \).

Since an edge of \( Q \) quotients at most \( \max_\omega \) edges of \( G \), \( f_E(\pi) \leq \max_\omega \cdot \left( cg(\varphi_Q) - \left[ \frac{\delta(\pi)}{2} \right] \right) ).

- The second term is straightforwardly bounded by \( |\omega(\pi)| \). However, from every vertex \( \omega' \) of \( \pi \) such that \( \varphi_G(\omega') \geq \varphi_G(\omega) \), there exist at least \( \delta_{\Pi,\varphi_Q}^+(\pi) \) edges of \( \omega(\pi) \) incident to vertices \( \omega'' \) with \( \varphi_G(\omega'') > \varphi_G(\omega) \), which must not be accounted for. Similarly, for each vertex \( \omega' \) of \( \pi \) such that \( \varphi_G(\omega') < \varphi_G(\omega) \), at least \( \delta_{\Pi,\varphi_Q}^-(\pi) \) edges must not be accounted for.

Therefore, \( f_C(\pi) \leq |\omega(\pi)| - |\pi| \cdot \min(\delta_{\Pi,\varphi_Q}^-(\pi), \delta_{\Pi,\varphi_Q}^+(\pi)) \).

- The third term is, by definition, equal to the congestion of \( G[\pi] \) with the restriction of \( \varphi_G \) to \( \pi \), i.e. \( cg(\varphi_G[\pi]) \).

By combining these three upper bounds, we obtain

\[ Cw(G) \leq \max_{\pi \in \Pi} \left( \left( cg(\varphi_Q) - \left[ \frac{\delta(\pi)}{2} \right] \right) \cdot \max_\omega + cg(\varphi_G[\pi]) \right. \]

\[ + \left. \left( |\omega(\pi)| - |\pi| \cdot \min(\delta_{\Pi,\varphi_Q}^-(\pi), \delta_{\Pi,\varphi_Q}^+(\pi)) \right) \right). \]

Since, in any block, the ordering of the vertices is continuous by definition, and independent from the orderings of other blocks, it is possible, once taken an ordering \( \varphi_Q \) which achieves the cutwidth of \( Q \), to obtain a compatible ordering \( \varphi_G \) in \( \Phi_G(G, \Pi, \varphi_Q) \) such that the restriction of \( \varphi_G \) to every block \( \pi \) in \( \Pi \) gives orderings which achieve the cutwidth of \( G[\pi] \). Combining this ordering with the above equation yields the claimed result.

This theorem leads to the more simple form stated in the next corollary.
**Corollary 6:** If $G$ is a connected graph and $Q$ is the quotient graph obtained from a partition $\Pi$ of $V(G)$ with at least two blocks, then

$$Cw(G) \leq Cw(Q) \Delta(G) \max_{\pi}.$$

**Proof:** If $G$ is connected and $\Pi$ is not restricted to a single block, then $\delta(Q) \geq 1$. For all blocks $\pi$, $Cw(G[\pi]) + |\omega(\pi)|$ is bounded by the number of edges having at least an end in $\pi$ and thus by $\Delta(G)|\pi|$, and $\max_{\omega} \leq \Delta(G) \max_{\pi}$. □

**Proposition 7:** For all graphs $G$,

$$Cw(G) \geq bis_e(G).$$

**Proof:** Let $\varphi_G$ be an ordering of $G$ which achieves the cutwidth of $G$, and let $E_d = \left\{ \{v', v''\} \in E(G)/\varphi_G(v') \leq \frac{n-1}{2} < \varphi_G(v'') \right\}$. By definition of the cutwidth, $|E_d| \leq Cw(G)$, and $E_d$ disconnects $G$ into two subgraphs $G_1$ and $G_2$ such that $|V(G_1)| = \left\lfloor \frac{n-1}{2} \right\rfloor$ and $|V(G_2)| = \left\lceil \frac{n-1}{2} \right\rceil$, so $bis_e(G) \leq |E_d| \leq Cw(G)$. □

### 4. Applications

**4.1. Binary de Bruijn graph**

We exclusively study the case of unoriented binary de Bruijn graphs, denoted $UB(2, k)$. The vertices $v$ of $UB(2, k)$ are words of $k$ letters taken in the $\{0, 1\}$ alphabet, and denoted $v = v_1 v_2 ... v_k$. There exists an edge between two distinct vertices $v$ and $v'$ if and only if the $(k-1)$ leftmost letters of one of the two vertices are equal to the $(k-1)$ rightmost letters of the other vertex, *i.e.* if $v' = v_2 v_3 ... v_k x$ or $v' = x v_1 ... v_{k-2} v_{k-1}$, with $x \in \{0, 1\}$.

We define the complement of $v$ as $\bar{v} = \bar{v}_1 \bar{v}_2 ... \bar{v}_k$, where $\bar{v}_i = 1 - v_i$; by extension, the complement of an edge is the edge whose ends are the complements of the ends of this edge. In reference to language theory, we may denote $x^k$ a letter $x$ repeated $k$ times within a word. We denote $\mathcal{H}(v)$ the Hamming weight of $v$, *i.e.* the number of “1”s in the word representation of $v$. If $v$ is a vertex to which is associated a word on a binary alphabet, we define as $[v]$ the value of this word taken as a binary number representation.
PROPOSITION 8:

\[ \text{Bd} \left( \text{UB}(2, k) \right) \leq \binom{k}{\left\lfloor \frac{k}{2} \right\rfloor}. \]

Proof: In order to obtain a quotient graph which is a path, we quotient \( \text{UB}(2, k) \) using the Hamming weight of its vertices, since the difference of the Hamming weight of any two vertices linked by an edge is at most 1 by definition. Therefore, we define the following configuration:

- \( \Pi = \{\pi_0, \pi_1, ..., \pi_k\} \), where \( \pi_q = \{v \in V(\text{UB}(2, k))/H(v) = q\} \);
- \( \varphi_Q(\pi_q) = q \);
- \( \varphi_G \), defined as:

  \( \varphi_G \), which is compatible with \( \varphi_Q \), amounts, in each block, to numbering the vertices \( v \) in decreasing order with respect to \( [v] \); it is therefore equivalent to the ordering defined by Harper for the hypercube \([15]\).

This ordering has several interesting properties:

- \( \text{UB}(2, k)/\Pi \) is isomorphic to the path graph \( P(k + 1) \).

- \( |\pi_q| = \binom{k}{q} \leq \left\lfloor \frac{k}{\frac{k}{2}} \right\rfloor \).

- \( (v \in \pi_q) \Leftrightarrow (\bar{v} \in \pi_{k-q}) \).
\[ \varphi_G (\overline{v}) = 2^k - 1 - \varphi_G (v). \]

Since \( UB (2, k) /_\Pi \) is isomorphic to the path graph \( P (k + 1) \), then \( Bd (UB (2, k) /_\Pi) = 1 \) and \( Bd (UB (2, k)) = \max (\varepsilon_i (\varphi_G), \varepsilon_e (\varphi_G)) \). Moreover, for each edge \( \{v', v''\} \) of \( E (\Pi) \) such that \( \varphi_G (v') < \varphi_G (v'') \), we have \( P_{\varphi_Q (\pi (v'))} = P_{\varphi_Q (\pi (v''))} = 1 \), and thus \( \varepsilon_e (\varphi_G) = \max_{\{v', v''\} \in E (\Pi)} (|\varphi_G (v'') - \varphi_G (v')|). \)

Let us consider all the edges \( \{v', v''\} \) in \( E (UB (2, k)) \), assuming without loss of generality that \( q' = \varphi_Q (\pi (v')) \leq \varphi_Q (\pi (v'')) = q'' \).

\begin{itemize}
  \item If \( q'' = q' \), then \( |\varphi_G (v') - \varphi_G (v'')| \leq |\pi_{q'}| - 1 \leq \left( \frac{k}{q'} \right) - 1 \), and then \( \varepsilon_i (\varphi_G) \leq \left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right) - 1. \)

  \item If \( q'' = q' + 1 \), then \( v'' \) is obtained from \( v \) by deletion of a letter “0” and insertion of a letter “1”.

  \begin{itemize}
    \item If \( v' = 0 m \) and \( v'' = 1 m \), let us consider vertex \( v'_\perp = 0^{k-q} 1^q = 0 m_\perp \) in \( \pi_{q'} \) and its neighbor \( v''_\perp = m_\perp 1 \) in \( \pi_{q''} \). By definition, for all \( m, [m_\perp] \leq [m] \), and \( ([0 m] - [0 m_\perp]) \leq ([m] - [m_\perp 1]) \). Since \( (\varphi_G (v'_\perp) - \varphi_G (v')) = ([0 m] - [0 m_\perp]) \) and \( (\varphi_G (v''_\perp) - \varphi_G (v'')) = ([m] - [m_\perp 1]), \) we have \( (\varphi_G (v'') - \varphi_G (v')) \leq (\varphi_G (v'_\perp) - \varphi_G (v'_\perp)) = |\pi_{q''}| \leq \left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right). \)

    \item If \( v' = 0 m \) and \( v'' = 1 m \), then \( \overline{v'} = \overline{m} 1 \) and \( \overline{v''} = 0 \overline{m} \), which amounts to considering the first case since \( |\varphi_G (v'') - \varphi_G (v')| = |\varphi_G (\overline{v''}) - \varphi_G (\overline{v'})| \).

  \end{itemize}

By the above, \( \varepsilon_e (\varphi_G) = \left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right). \)

By Theorem 1, taking the maximum of \( \varepsilon_i (\varphi_G) \) and \( \varepsilon_e (\varphi_G) \) yields the claimed result. \( \square \)

*Remark:* \( \frac{2^k - 1}{k} \leq Bd (UB (2, k)). \)

*Proof:* This derives straightforwardly from Theorem 4, with \( |V (UB (2, k))| = 2^k \) and \( \text{diam} (UB (2, k)) = k \).
PROPOSITION 9:

\[
\frac{1}{2} \cdot \frac{2^k}{k} \leq Cw(UB(2, k)) \leq 2 \left( \frac{k}{\left\lfloor \frac{k}{2} \right\rfloor} \right) + 2.
\]

Proof: To prove the upper bound of this proposition, we use a configuration based on the same \( \Pi \) and \( \varphi_Q \) as above. Since \( UB(2, k)_{/\Pi} \) is isomorphic to \( P(k + 1) \), \( Cw(UB(2, k)_{/\Pi}) = 1 \) and \( \delta(UB(2, k)_{/\Pi}) = 1 \). Moreover, \(|\omega(\pi_q)| \leq 2|\pi_q|\), as from each vertex in any \( \pi \) exit at most 2 edges to other \( \pi \)'s.

The edges \( \{v', v''\} \) internal to any block \( \pi \) link vertices of same weight, so if \( v' = v_1 v_2 \ldots v_k \), then \( v'' = v_k v_1 \ldots v_{k-1} \) or \( v'' = v_2 \ldots v_k v_1 \). Therefore, the connected components of \( UB(2, k)[\pi] \) are single vertices, single edges, or cycles (named Etzion cycles) and, for all \( \pi \), \( Cw(UB(2, k)[\pi]) = 2 \).

For \( k \geq 3 \) and for any block \( \pi_q \), it is possible to find a vertex not linked to either a vertex of \( \pi_{q-1} \) or \( \pi_{q+1} \) (i.e. a vertex whose first and last letters are identical) so, for all \( \pi \), \( \min(\delta^-_{\Pi, \varphi_Q}(\pi), \delta^+_{\Pi, \varphi_Q}(\pi)) = 0 \).

According to the above, Theorem 5 yields

\[
Cw(UB(2, k)) \leq \left( Cw(UB(2, k)_{/\Pi}) - \left\lfloor \frac{\delta(UB(2, k)_{/\Pi})}{2} \right\rfloor \right) \cdot \max_\omega \\
+ \max_{\pi \in \Pi} \left( Cw(UB(2, k)[\pi]) + |\omega(\pi)| \right) \\
- |\pi| \cdot \min(\delta^-_{\Pi, \varphi_Q}(\pi), \delta^+_{\Pi, \varphi_Q}(\pi)) \\
\leq \max_{\pi \in \Pi} \left( Cw(UB(2, k)[\pi]) + |\omega(\pi)| \right) \\
\leq 2 + 2 \left( \frac{k}{\left\lfloor \frac{k}{2} \right\rfloor} \right).
\]

It is easy to see that \( Cw(UB(2, 1)) = 1 \) and \( Cw(UB(2, 2)) = 3 \), which extends the upper bound to all \( k \geq 1 \).

For the lower bound, it is known from [19, p. 480] that \( bis_e(SE(k)) \geq \frac{2^{k-1}}{k} \). Since \( SE(k) \) is a spanning subgraph of \( UB(2, k) \),

\[
Cw(UB(2, k)) \geq bis_e(UB(2, k)) \geq bis_e(SE(k)) \geq \frac{1}{2} \cdot \frac{2^k}{k}. \quad \Box
\]
Remark: It is known that, for all \( k > 1 \),
\[
\left( \left\lfloor \frac{k}{2} \right\rfloor \right) < 2^k \frac{2}{k^2}.
\]
Hence, our upper bounds of the bandwidth and the cutwidth of the binary de Bruijn graphs are both in \( O \left( \frac{2^k}{\sqrt{k}} \right) \).

### 4.2. Butterfly graph

The vertices of \( BF(k) \) are pairs of integer numbers \((l; m)\), where \( 0 \leq l < k \) and \( 0 \leq m < 2^k \), and each vertex \((l; m)\) of \( BF(k) \) is linked to vertices \(((l - 1) \mod k; m), ((l - 1) \mod k; m \oplus 2^k), ((l + 1) \mod k; m),\) and \(((l + 1) \mod k; m \oplus 2^{(l+1) \mod k})\). Vertices \((l; m)\) of same \( l \) value are said to belong to the same level \( l \).

**Proposition 10:**

\[
Bd(BF(k)) \leq 9 \cdot 2^{k-2}.
\]

**Proof:** The first part of the proof is devoted to the definition of a suitable configuration. The \( \varphi_G \) numbering that we obtain is then studied to achieve a better upper bound.

A natural way to partition Butterfly graphs is to quotient them by levels, as illustrated in Figure 3. We will thus define \( \Pi, \varphi_Q, \) and \( \varphi_G \) as follows:

![Figure 3. – BF (3) and its partition by levels.](image-url)
\[ \Pi = \{ \pi_0, \pi_1, \ldots, \pi_{k-1} \}, \text{ where } \pi_q = \{(l; m) \in V(BF(k))/l = q \} \]

\[ \varphi_Q(\pi_q) = \begin{cases} 
2(q + 1) & \text{if } 0 \leq q < \left\lfloor \frac{k - 1}{2} \right\rfloor \\
2(k - (q + 1)) - 1 & \text{if } \left\lfloor \frac{k - 1}{2} \right\rfloor \leq q < k - 1 \\
0 & \text{if } q = k - 1 
\end{cases} \]

\[ \varphi_G(l; m) = \varphi_Q(\pi_l) \cdot 2^k + m. \]

By definition of \( \Pi \), it is clear that, for all \( q \), \( |\pi_q| = 2^k \), and that \( \varphi_G \) is compatible with \( \varphi_Q \). Since no edge links vertices belonging to the same block, \( \epsilon_1(\varphi_G) = 0 \). By definition, \( BF(k)_{/\Pi} \) is isomorphic to the cycle graph \( C(k) \), so \( Bd(BF(k)_{/\Pi}) = 2 \). Moreover, the above definition of \( \varphi_Q \) is an ordering which achieves this bandwidth on the cycle.

For all \( v = (l; m) \in V(BF(k)) \), \( v \) is linked to \( v' = ((l + 1) \mod k; m) \) and to \( v'' = ((l + 1) \mod k; m \oplus 2^{(l+1) \mod k}) \). Thus,

\[ \epsilon_e(\varphi_G) = 2^k + 2^{k-1}. \]

However, with the chosen \( \varphi_Q \), all the edges which maximize \( \epsilon_e(\varphi_G) \) link block \( \pi_{k-2} \) to block \( \pi_{k-1} \), and \( (\varphi_Q(\pi_{k-2}) - \varphi_Q(\pi_{k-1})) \) is by construction always equal to 1 (and not to 2, which is the bandwidth of \( Q \)). Therefore, the maximum dilation of the external edges is in fact less than what would be obtained by using Theorem 1. As a matter of fact,

\[ Bd(BF(k)) \leq \max(2^k \cdot |\varphi_Q(\pi_{k-1}) - \varphi_Q(\pi_{k-2})| + 2^{k-1}, 2^k \cdot |\varphi_Q(\pi_{k-2}) - \varphi_Q(\pi_{k-3})| + 2^{k-2}) \]

\[ \leq \max(1 \cdot 2^k + 2^{k-1}, 2 \cdot 2^k + 2^{k-2}) \]

\[ \leq 9 \cdot 2^{k-2}. \]

**Remark:**

\[ \left\lceil \frac{8}{3} \cdot 2^{k-2} \right\rceil \leq Bd(BF(k)) \]

**Proof:** As we know by [19] that \( \text{diam}(BF(k)) = \left\lfloor \frac{3}{2} k \right\rceil \), we have by Theorem 4

\[ Bd(BF(k)) \geq \left\lceil \frac{k \cdot 2^k - 1}{\left\lfloor \frac{3}{2} k \right\rceil} \right\rceil \geq \left\lceil \frac{k \cdot 2^k - 1}{\left\lfloor \frac{3}{2} k \right\rceil} \right\rceil \]

\[ = \left\lceil \frac{8}{3} \cdot 2^{k-2} - \frac{2}{3} k \right\rceil \geq \left\lceil \frac{8}{3} \cdot 2^{k-2} \right\rceil, \]

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as the contents of the ceiling operator are never integer. □

**Proposition 11:**

\[
\left[ \frac{1}{5} \cdot 2^{k+1} \right] \leq C_w(BF(k)) \leq 3 \cdot 2^k.
\]

**Proof:** Let us begin with the upper bound. An interesting characteristic of Butterfly graphs is that their construction is recursive. As a matter of fact, let us define the BF graphs as BF graphs without their wrap-around edges (i.e., edges linking level \( k - 1 \) to level 0). It is easy to see that BF \( (k) \) can be obtained by putting two copies of BF \( (k - 1) \) atop a level of \( 2^k \) vertices and linking them so that the added level becomes level \( k \), as illustrated in figure 4. This decomposition is straightforward for BF \( (k) \), as illustrated in figure 5.

![Figure 4. BF (3) and its partition: two BF (2) and a level of vertices.](image)

![Figure 5. Recursive structure of BF graphs.](image)
To prove the upper bound, we use this two-step recursive construction of BF (k). The first step, involving the decomposition of BF (k), leads to the following \( \Pi \) and \( \varphi_Q \):

- \( \Pi = \{ \pi_0, \pi_1, \pi_2 \} \), where
  \[
  \begin{cases}
  \pi_0 = \{(l; m) \in V(BF(k))/l < k - 1, m < 2^{k-1}\} \\
  \pi_1 = \{(l; m) \in V(BF(k))/l = k - 1\} \\
  \pi_2 = \{(l; m) \in V(BF(k))/l < k - 1, m \geq 2^{k-1}\}
  \end{cases}
  \]

- \( \varphi_Q(\pi_q) = q \).

By definition, \( \pi_0 \) and \( \pi_2 \) are isomorphic to \( BF(k-1) \) graphs, and the quotient graph \( BF(k)/\Pi \) is a path, so \( Cw(BF(k)/\Pi) = 1 \), \( \delta(BF(k)/\Pi) = 1 \), and \( BF(k)[\pi_1] \) is a set of independent vertices, so \( Cw(BF(k)[\pi_1]) = 0 \). This structure yields the following values:

- \( \left| \pi_0 \right| = (k - 1)2^{k-1} \), \( \left| \omega(\pi_0) \right| = 4 \cdot 2^{k-1}, Cw(BF(k)[\pi_0]) = Cw(BF(k-1)) \),
- \( \left| \pi_1 \right| = 2^k \), \( \left| \omega(\pi_1) \right| = 4 \cdot 2^k, Cw(BF(k)[\pi_1]) = 0 \),
- \( \left| \pi_2 \right| = (k - 1)2^{k-1} \), \( \left| \omega(\pi_2) \right| = 4 \cdot 2^{k-1}, Cw(BF(k)[\pi_2]) = Cw(BF(k-1)) \),

\[
\begin{align*}
\delta_{\Pi, \varphi_Q}(\pi_0) &= 0, & \delta^+_{\Pi, \varphi_Q}(\pi_0) &= 0, \\
\delta_{\Pi, \varphi_Q}(\pi_1) &= 1, & \delta^+_{\Pi, \varphi_Q}(\pi_1) &= 1, \\
\delta_{\Pi, \varphi_Q}(\pi_2) &= 0, & \delta^+_{\Pi, \varphi_Q}(\pi_2) &= 0.
\end{align*}
\]

Using Theorem 5, we obtain

\[
Cw(BF(k)) \leq \max_{\pi \in \{\pi_0, \pi_1, \pi_2\}} (Cw(BF(k)[\pi]) + |\omega(\pi)| - |\pi| \cdot \min(\delta_{\Pi, \varphi_Q}(\pi), \delta^+_{\Pi, \varphi_Q}(\pi))) + \min(\delta_{\Pi, \varphi_Q}(\pi), \delta^+_{\Pi, \varphi_Q}(\pi)))
\]

\[
\leq \max(Cw(BF(k-1)) + 2^{k+1}, 3 \cdot 2^k).
\] (1)

In the second step, an analogous study on \( BF(k) \) based on the same partition leads to:

\[
Cw(BF(k)) \leq \max(Cw(BF(k-1)) + 2^k, 2^k) \leq Cw(BF(k-1)) + 2^k.
\] (2)
Since $\overline{BF}(2)$ is isomorphic to two $C(4)$, $Cw(\overline{BF}(2)) = 2$. By injecting this value into (2), and studying the induced recurrence, we obtain

$$Cw(\overline{BF}(k)) \leq 2^{k+1} - 6$$

which, injected itself into (1), yields the claimed result.

Let us now prove the lower bound. The Cube-Connected Cycles graph of dimension $k$, denoted $CCC(k)$, is the graph obtained by replacing the vertices of a $k$-dimensional hypercube by $k$-node cycles [26]. In [27] has been given a lower bound of the edge-bisection of a graph with respect to the congestion of the embedding of the graph into the complete graph with double edges. In particular, for the CCC graph, this approach yields

$$\text{bis}_e (CCC(k)) \geq \frac{1}{5} \cdot 2^{k+1}.$$}

Since it has been proven in [8] that $CCC(k)$ is a spanning subgraph of $BF(k)$, $Cw(CCC(k)) \leq Cw(BF(k))$, and by Proposition 7

$$Cw(BF(k)) \geq \left\lfloor \frac{1}{5} \cdot 2^{k+1} \right\rfloor. \quad \square$$

### 4.3. Summary

Using the same techniques as above, our method has been successfully applied to other classes of graphs [2]. All these results are summarized in the following table.

**Remark:** The upper bound that we have found for $Cw(SE(k))$ is in $O\left(\frac{2^k}{\sqrt{k}}\right)$. However, by combining the constant-congestion embedding of $SE(k)$ into $BF(k + 2 - \log(k))$ described in [16, p. 237] with the embedding that we use to compute the upper bound of $Cw(BF)$, one can prove that $Cw(SE(k))$ is in $O\left(\frac{2^k}{k}\right)$. Since the lower bound that we give is also of this order of magnitude, then $Cw(SE(k))$ is in $\Theta\left(\frac{2^k}{k}\right)$.

### 5. Conclusion

We have presented in this paper a method for determining upper bounds of the bandwidth and the cutwidth of a graph from the bandwidth and the vol. 29, n° 6, 1995
Graph $\Theta \left( \left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right) \right)$ [17] $\left[ \frac{1}{3} 2^{k+1} \right]$ [3, 4, 21]
UB(2, $k$) $\frac{2^{k-1}}{k}$ $\left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right)$ $\frac{1}{2} \cdot \frac{2^k}{k}$ $2 \left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right) + 2$
SE($k$) $\frac{2^{k-1}}{2k-1}$ $2 \left( \frac{k-2}{\left\lceil \frac{k-2}{2} \right\rceil} \right)$ $\frac{1}{2} \cdot \frac{2^k}{k}$ $\uparrow$ $\left( \frac{k}{\left\lceil \frac{k}{2} \right\rceil} \right) + 2 \uparrow$
FFT($k$) $2^{k-1}$ $3 \cdot 2^{k-1}$ $\left[ \frac{1}{5} \cdot 2^{k+1} \right]$ $2^{k+1} - 2$
BF($k$) $\left[ \frac{8}{3} \cdot 2^{k-2} \right]$ $9 \cdot 2^{k-2}$ $\left[ \frac{1}{5} \cdot 2^{k+1} \right]$ $3 \cdot 2^k$
CCC($k$) $\left[ \frac{8}{5} \cdot 2^{k-2} \right]$ $9 \cdot 2^{k-2}$ $\left[ \frac{1}{5} \cdot 2^{k+1} \right]$ $2^{k+1} + 1$

† $Cw(SE(k))$ is in fact $\Theta \left( \frac{2^k}{k} \right)$ (see remark above).

The quotienting method is well suited for graphs which can be quotiented, possibly recursively, into graphs the bandwidth or cutwidth of which are known or for which accurate upper bounds are available. However, building suitable quotient graphs is not easy, since it inherits the complexity of the problem. Quotientings must be simple so that their parameters can be easily computed, and must be judiciously chosen in order to capture in the quotient graph the topological properties of the studied graphs.

When the graph structure is well-known, it is often worth fine-tuning the ordering of the vertices induced by the quotienting, in order to directly compute more accurate upper bounds. In that way, our results can be seen as a methodological framework for finding good orderings rather than directly applicable theorems.

Applications of the method to the binary de Bruijn and Butterfly graphs have been exposed in the paper. Studies of other classes of interconnection networks (Shuffle-Exchange, FFT, CCC), and generalizations of these results to $d$-ary graphs have also been carried out in related articles [1, 2, 23, 24].
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