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### On the minimum of the unit lattice.

## PAR VOLKER KESSLER

#### 1. Introduction.

Computations in lattices often require a lower bound for the minimum of the lattice, both for practical purposes and for a theoretical analysis of the algorithms, e.g. [1] and [2].

In this paper we recall two results of Dobrowolski [3] and Smyth [5] in order to get such a bound for the unit lattice.

### 2. Lower bound.

Let K be a finite extension of  $\mathbb{Q}$  of degree n with maximal order R. For  $1 \leq i \leq n$  we denote by

$$K \to K^{(i)} \subset \mathbb{C}, \ \alpha \to \alpha^{(i)}$$

the n different embeddings of K into the field  $\mathbb{C}$  of complex numbers. The first  $r_1$  of those embeddings are real, the last  $2r_2$  embeddings are non-real and numbered such that the  $(r_1 + r_2 + i)$ th embedding is the complex-conjugation of the  $(r_1 + i)$ th embedding. Then the logarithmic map is given by

$$\operatorname{Log}: K^* \to \mathbb{R}^r, \quad \operatorname{Log}(\alpha) := (c_1 \log |\alpha^{(1)}|, \cdots, c_r \log |\alpha^{(r)}|)$$

with the unit rank  $r = r_1 + r_2 - 1$  and

$$c_i = \begin{cases} 1 & \text{for } 1 \le i \le r_1 \\ 2 & \text{for } r_1 + 1 \le i \le r + 1. \end{cases}$$

The kernel of Log consists exactly of the roots of the unity lying in K. We define the minimum  $\lambda(L)$  of the unit lattice  $L := \text{Log}(R^*)$  by

$$\lambda(L) = \min\{ ||v|| | v \in L \setminus \{0\} \}$$

where | | | denotes the Euclidean norm.

THEOREM: A lower bound for the minimum  $\lambda(L)$  is given by (1)

$$\lambda(L) > \mu(K) := \sqrt{\frac{2}{r+1}} \left( \frac{1}{1200} \left( \frac{\log \log n}{\log n} \right)^3 - \frac{1}{2880000} \left( \frac{\log \log n}{\log n} \right)^6 \right)$$

which is "a bit" larger than

$$\frac{1}{\sqrt{r+1}} \frac{1}{1000} \left( \frac{\log \log n}{\log n} \right)^3.$$

Thus the inverse  $1/\lambda(L)$  is of the magnitude  $0(n^{1/2+\epsilon})$  for every  $\epsilon > 0$ .

PROOF. Let  $\epsilon \in R^*$  be a unit of degree m over  $\mathbb{Q}$ , which is no root of unity. Without loss of generality we can assume that m=n, because if  $\|\text{Log }\epsilon\|$  is larger than  $\mu(K')$  for a subfield K' of K it is also larger than  $\mu(K)$ .

We are interested in two subsets of the conjugates  $\epsilon^{(1)}, \dots, \epsilon^{(n)}$ 

$$S := \{1 \le i \le r + 1 \mid |\epsilon^{(i)}| > 1\}$$
$$T := \{1 \le i \le r + 1 \mid |\epsilon^{(i)}| < 1\}.$$

Since  $\epsilon$  is no root of unity S is non-empty and therefore T cannot be empty because of  $N(\epsilon) = 1$ .

We call  $\epsilon$  reciprocal if  $\epsilon$  is conjugate to  $\epsilon^{-1}$ , i.e. its minimal polynomial  $f(X) = X^n + a_{n-1}X^{n-1} + \cdots + a_0$  satisfies

$$f(X) = X^n f(\frac{1}{X}) = a_0 X^n + a_1 X^{n-1} + \dots + a_{n-1} X + 1.$$

If  $\epsilon$  is non-reciprocal we know from the theorem of [5] that

$$\prod_{i \in S} |\epsilon^{(i)}|^{c_i} \ge \theta$$

where  $\theta$  is the real root of  $X^3 - X - 1$ , i.e.  $\theta \approx 1.3247$ . Thus

(2) 
$$\sum_{i \in S} c_i \log |\epsilon^{(i)}| \ge \log \theta \approx 0.281$$

But from  $N(\epsilon) = 1$  it follows

(3) 
$$\sum_{i \in S} c_i \log |\epsilon^{(i)}| = -\sum_{i \in T} c_i \log |\epsilon^{(i)}|.$$

The value  $c_{r+1} \log |\epsilon^{(r+1)}|$  does not occur in the norm of  $\operatorname{Log}(\epsilon)$ . But as a consequence of (3) it does not matter if r+1 lies in S or in T and so we can assume without restriction that  $r+1 \notin S$ . Thus

$$\begin{aligned} ||\operatorname{Log}(\epsilon)|| &\geq \sqrt{\sum_{i \in S} (c_i |\log |\epsilon^{(i)}|)^2} \\ &\geq r^{-1/2} \sum_{i \in S} (c_i |\log |\epsilon^{(i)}|) \geq r^{-1/2} \log |\theta| > \mu(K). \end{aligned}$$

(The second inequality follows from the well known norm equivalence between 1-norm and Euclidean norm.)

For reciprocal  $\epsilon$  we know by Theorem 1 of [3]:

(4) 
$$\prod_{i \in S} |\epsilon^{(i)}|^{c_i} > 1 + \frac{1}{1200} \left( \frac{\log \log n}{\log n} \right)^3.$$

We now use the Taylor series of the logarithm (|y| < 1):

(5) 
$$\log(1+y) = y - \frac{y^2}{2} + \frac{y^3}{3} \mp \dots > y - \frac{y^2}{2}.$$

The inequality follows directly from Lagrange's representation of the residue. Applying (5) to (4) yields

$$\sum_{i \in S} c_i \log |\epsilon^{(i)}| > \frac{1}{1200} (\frac{\log \log n}{\log n})^3 - \frac{1}{2880000} (\frac{\log \log n}{\log n})^6.$$

Since  $\epsilon$  is reciprocal the inverses of the conjugates of  $\epsilon$  are also conjugate to  $\epsilon$ . This implies that the numbers of conjugates outside the unit circle equals the number of conjugates inside the unit circle, i.e

$$\#S = \#T \le \frac{r+1}{2} \le \frac{n}{2}.$$

Again by (3) we can assume that  $r+1 \notin S$ 

$$||\operatorname{Log}(\epsilon)|| \ge \sqrt{\sum_{i \in S} (c_i \log |\epsilon^{(i)}|)^2} \ge \sqrt{\frac{2}{r+1}} \sum_{i \in S} c_i \log |\epsilon^{(i)}|$$

$$> \sqrt{\frac{2}{r+1}} \left( \frac{1}{1200} (\frac{\log \log n}{\log n})^3 - \frac{1}{2880000} (\frac{\log \log n}{\log n})^6 \right) = \mu(K)$$

which is larger than

$$\sqrt{\frac{2}{r+1}}(\frac{1}{1200} - \frac{1}{2880000})(\frac{\log \log n}{\log n})^3.$$

Because of  $\sqrt{2}(\frac{1}{1200} - \frac{1}{2880000}) \approx 0.001178$  we thus proved the lower bound.

REMARK. If the conjecture of Schinzel and Zassenhaus [5] is correct the term  $(\frac{\log \log n}{\log n})^3$  can be substituted by a constant independent of n. This bound would be provable the best one (up to constants).

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