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Number Theory / *Théorie des nombres*

The subword complexity of polynomial subsequences of the Thue–Morse sequence

La complexité de facteurs des sous-suites polynomiales de la suite de Thue–Morse

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Abstract. Let $\mathbf{t} = (t(n))_{n \geq 0}$ be the Thue–Morse sequence in $\{0, 1\}$. J.-P. Allouche and J. Shallit asked in 2003 whether the subword complexity of the subsequence $(t(n^2))_{n \geq 0}$ attains the maximal value. This problem was solved positively by Y. Moshe in 2007. Indeed Y. Moshe had shown that for all $H \in \mathbb{Q}[T]$ with $H(\mathbb{N}) \subseteq \mathbb{N}$ and $\deg H = 2$, all the subsequences $(t(H(n)))_{n \geq 0}$ attain the maximal subword complexity. Then he asked whether the same result holds for $\deg H \geq 3$. In this work, we shall give a positive answer to the above problem.

Résumé. Soit $\mathbf{t} = (t(n))_{n \geq 0}$ la suite de Thue–Morse en $\{0, 1\}$. J.-P. Allouche et J. Shallit demandaient en 2003 si la complexité de facteurs de la sous-suite $(t(n^2))_{n \geq 0}$ atteint la maximale. Le problème était résolu positivement par Y. Moshe en 2007. En fait, Y. Moshe avait démontré que pour tout $H \in \mathbb{Q}[T]$ avec $H(\mathbb{N}) \subseteq \mathbb{N}$ et $\deg H = 2$, toutes les sous-suites $(t(H(n)))_{n \geq 0}$ atteignent la complexité maximale. Ensuite il demandait si le résultat est aussi valable pour $\deg H \geq 3$. Dans ce travail, nous allons donner une réponse positive au problème précédent.

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Soient Σ un alphabet fini et $\mathbf{u} = (u(n))_{n \geq 0}$ une suite sur Σ . Pour tout entier $m \geq 1$, désignons par $P_{\mathbf{u}}(m)$ le nombre de facteurs différents de longueur m dans \mathbf{u} , et appelons la fonction $P_{\mathbf{u}}$ la complexité de facteurs de \mathbf{u} . Ainsi $P_{\mathbf{u}}(m) \leq |\Sigma|^m$, où $|\Sigma|$ désigne le nombre d'éléments dans Σ .

Soit $\mathbf{t} = (t(n))_{n \geq 0}$ la suite de Thue–Morse en $\{0, 1\}$. Elle est définie par $t(n) = s_2(n) \pmod{2}$, pour tout entier $n \geq 0$, où $s_2(n)$ désigne le nombre de 1's dans la représentation binaire de n . La complexité de facteurs de \mathbf{t} est compliquée mais déjà connue (voir S. Brlek [4], A. de Luca

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et S. Varricchio [8], et S. V. Avgustinovich [3]). J.-P. Allouche et J. Shallit demandaient dans [2] si $P_{\mathbf{u}}(m) = 2^m$ pour tout $m \geq 1$, où $\mathbf{u} = (t(n^2))_{n \geq 0}$. Le problème était résolu positivement par Y. Moshe [9]. Plus généralement, il a obtenu la même conclusion pour toutes les sous-suites $(t(H(n)))_{n \geq 0}$, avec $H \in \mathbb{Q}[T]$ tel que $H(\mathbb{N}) \subseteq \mathbb{N}$ et $\deg H = 2$, et il demandait ensuite si le résultat persiste encore pour $\deg H \geq 3$.

Dans la suite, nous allons donner une réponse positive au problème précédent de Y. Moshe.

Theorem. *Soit $H \in \mathbb{Q}[T]$ tel que $H(\mathbb{N}) \subseteq \mathbb{N}$ et $\deg H \geq 2$. Alors $P_{\mathbf{u}}(m) = 2^m$ pour tout entier $m \geq 1$, où $\mathbf{u} = (u(n))_{n \geq 0} = (t(H(n)))_{n \geq 0}$.*

1. Introduction

Let Σ be a finite alphabet and $\mathbf{u} = (u(n))_{n \geq 0}$ a sequence over Σ . For all integers $m \geq 1$, let $P_{\mathbf{u}}(m)$ be the number of different subwords in \mathbf{u} with length m , and we call the function $P_{\mathbf{u}}$ the subword complexity of \mathbf{u} . So $P_{\mathbf{u}}(m) \leq |\Sigma|^m$, where $|\Sigma|$ denotes the number of elements in Σ . Note that automatic sequences have relatively low complexity $O(m)$, and random sequences have high complexity (see e.g. [2]).

Let $\mathbf{t} = (t(n))_{n \geq 0}$ be the Thue–Morse sequence in $0, 1$, i.e., $t(n) = s_2(n) \pmod{2}$, where $s_2(n)$ is the number of 1’s in the binary representation of n . The subword complexity of \mathbf{t} is complicated but already known (see S. Brlek [4], A. de Luca and S. Varricchio [8], and S. V. Avgustinovich [3]). It is well known that \mathbf{t} is 2-automatic, and J.-P. Allouche showed in [1] that $(t(H(n)))_{n \geq 0}$ is not 2-automatic, if $H \in \mathbb{Q}[T]$ with $H(\mathbb{N}) \subseteq \mathbb{N}$ and $\deg H \geq 2$. So $(t(n^2))_{n \geq 0}$ is not 2-automatic, and then J.-P. Allouche and J. Shallit asked in [2] whether $P_{\mathbf{u}}(m) = 2^m$ for all integers $m \geq 1$, where $\mathbf{u} = (t(n^2))_{n \geq 0}$. This problem was solved positively by Y. Moshe [9]. More generally, he has obtained the same conclusion for all $H \in \mathbb{Q}[T]$ with $H(\mathbb{N}) \subseteq \mathbb{N}$ and $\deg H = 2$, and then asked whether it holds also for $\deg H \geq 3$. Below we shall give a positive answer to this problem. For related works, see for example [5–7, 10, 11] and references therein.

Theorem 1. *Let $H \in \mathbb{Q}[T]$ such that $H(\mathbb{N}) \subseteq \mathbb{N}$ and $\deg H \geq 2$. Let $\mathbf{u} = (u(n))_{n \geq 0} = (t(H(n)))_{n \geq 0}$. Then $P_{\mathbf{u}}(m) = 2^m$ for all integers $m \geq 1$.*

Below let v_2 be the 2-adic valuation, and put $\mu(n) = 2^{v_2(n)}$ for all integers $n \geq 1$. For $\mathbf{e} = (e_j)_{1 \leq j \leq n}$, $\mathbf{f} = (f_j)_{1 \leq j \leq n} \in \mathbb{N}^n$, we say $\mathbf{e} < \mathbf{f}$ if $\exists k \in \mathbb{N}(1 \leq k \leq n)$ such that $e_j = f_j$ ($k < j \leq n$) and $e_k < f_k$, and call it the colexicographic order. Finally put

$$|\mathbf{e}| := \sum_{k=1}^n e_k, \quad \text{and} \quad C_{\mathbf{e}} := \binom{|\mathbf{e}|}{e_1, e_2, \dots, e_n} = |\mathbf{e}|! / \prod_{k=1}^n e_k!.$$

2. Some preliminary lemmas

Lemma 2. *Let $\ell \geq 1$ be an integer, and $\mathbf{e} = (e_j)_{1 \leq j \leq 2^\ell} \in \mathbb{N}^{2^\ell}$ with $1 \leq |\mathbf{e}| < 2^\ell$. Let $J_{\mathbf{e}}$ be the set of $\mathbf{f} \in \mathbb{N}^{2^\ell}$ such that $|\mathbf{f}| = |\mathbf{e}|$ and $C_{\mathbf{f}} = C_{\mathbf{e}}$. Then $\text{Card}(J_{\mathbf{e}})$ is even.*

Lemma 3. *Let $d \geq 0$, $\ell \geq 1$ be integers such that $d \leq 2^\ell - 1$. Let Y_d be the set of*

$$\mathbf{e} = (e_j)_{2 \leq j \leq 2^\ell} \in \mathbb{N}^{2^\ell - 1} \quad \text{with} \quad \sum_{j=2}^{2^\ell} e_j = d,$$

and $0 \leq e_j \leq 2(2 \leq j \leq 2^\ell)$. Then $y_d := \text{Card}(Y_d)$ is odd if and only if $d \equiv 0, 1 \pmod{3}$.

Lemma 4. Let $d \geq 0, \ell \geq 2$ be integers such that $d \leq 2^\ell - 2$. Let G_d be the set of

$$e = (e_j)_{3 \leq j \leq 2^\ell} \in \mathbb{N}^{2^\ell - 2} \quad \text{with} \quad \sum_{j=3}^{2^\ell} e_j = d,$$

and $0 \leq e_j \leq 2 (3 \leq j \leq 2^\ell)$. Then $g_d := \text{Card}(G_d)$ is odd if and only if $d \equiv 0, 2 \pmod{6}$.

Lemma 5. Let $d \geq 2, m \geq 1$ be integers, $H_j \in \mathbb{Z}[T]$ with $H_j(\mathbb{N}) \subseteq \mathbb{N}$ and

$$H_j = \sum_{k=0}^d a_k^{(j)} T^k \quad (1 \leq j \leq m).$$

Then $\exists N \geq 1$ such that for all $n > N, \exists A_n \geq 1$ such that $t(H_j(n + A_n)) = t(H_j(n)) (1 \leq j \leq m)$.

Lemma 6. Let $d \geq 2, m \geq 1$ be integers, $H_j \in \mathbb{Z}[T]$ with $H_j(\mathbb{N}) \subseteq \mathbb{N}$ and

$$H_j = \sum_{k=0}^d a_k^{(j)} T^k \quad (1 \leq j \leq m)$$

such that $a_d^{(j)} = a_d > 0$ and $a_{d-1}^{(1)} < a_{d-1}^{(2)} < \dots < a_{d-1}^{(m)}$. Then $\exists N \geq 1$ such that for all integers $n > N, \exists A_n \geq 1$ such that $t(H_j(n + A_n)) = t(H_j(n)) (1 \leq j < m)$ and $t(H_m(n + A_n)) = t(H_m(n)) + 1$.

Proof of Theorem 1. Write $H = \sum_{k=0}^d a_k T^k$ with $a_k \in \mathbb{Q}$, and $a_d \neq 0$. Take $Q \geq 2$ an integer such that $Qa_k \in \mathbb{Z}$ for $0 \leq k \leq d$. For $j \geq 1$, put $H_j(T) = H(QT + j)$. Then $H_j \in \mathbb{Z}[T]$, and we need to show that for all $b_j \in \mathbb{Z}/2\mathbb{Z} (1 \leq j \leq m)$, we can find infinitely many integers n such that $t(H_j(n)) = b_j (1 \leq j \leq m)$.

By induction on m . If $m = 1$, by Lemma 6, $\exists N_1 \geq 1$ such that for all integers $n > N_1$, we can find $A_n \geq 1$ such that $t(H_1(n + A_n)) = t(H_1(n)) + 1$, hence $t(H_1(n)) = b_1$ or $t(H_1(n + A_n)) = b_1$.

Now assume that the result holds for $m - 1$ with $m \geq 2$. Then there are infinitely many integers n such that $t(H_j(n)) = b_j (1 \leq j < m)$. If $b_m = t(H_m(n))$, then the desired result holds. Otherwise $b_m = t(H_m(n)) + 1$, and by Lemma 6, $\exists N_m \geq 1$ such that for all integers $n > N_m, \exists A_n \geq 1$ such that $t(H_j(n + A_n)) = t(H_j(n)) (1 \leq j < m)$ and $t(H_m(n + A_n)) = t(H_m(n)) + 1$. So the desired result holds. □

3. Proofs of lemmas

Proof of Lemma 2. Let S_{2^ℓ} be the symmetric group on 2^ℓ letters. For $\sigma \in S_{2^\ell}$ and $\mathbf{f} = (f_j)_{1 \leq j \leq 2^\ell} \in J_e$, put $\mathbf{f}_\sigma = (f_{\sigma(j)})_{1 \leq j \leq 2^\ell} \in J_e, J_e(\mathbf{f}) = \{\mathbf{g} \in J_e : \mathbf{g}_\sigma = \mathbf{f}\}$, and it suffices to show that $\text{Card}(J_e(\mathbf{f}))$ is even. Assume that \mathbf{f} has r different values, each with multiplicity $s_i (1 \leq i \leq r)$. Then

$$\sum_{1 \leq j \leq r} s_j = 2^\ell, \quad \text{and} \quad \text{Card}(J_e(\mathbf{f})) = \binom{2^\ell}{s_1, s_2, \dots, s_r}.$$

Note that $s_i < 2^\ell (1 \leq i \leq r)$, otherwise $f_j = f_1 (1 \leq j \leq 2^\ell)$, and then $2^\ell > |\mathbf{f}| = 2^\ell f_1$. Finally we obtain

$$v_2(\text{Card}(J_e(\mathbf{f}))) = \sum_{j=1}^{\ell} \left(\left\lfloor \frac{2^\ell}{2^j} \right\rfloor - \sum_{k=1}^r \left\lfloor \frac{s_k}{2^j} \right\rfloor \right) \geq \left\lfloor \frac{2^\ell}{2^\ell} \right\rfloor - \sum_{k=1}^r \left\lfloor \frac{s_k}{2^\ell} \right\rfloor = 1. \quad \square$$

Proof of Lemma 3. Write $(1 + x + x^2)^{2^\ell - 1} = \sum_{i=0}^{2(2^\ell - 1)} b_i x^i$, and set $\beta_k = \sum_{i=0}^k b_i (1 \leq k < 2^\ell)$. Then

$$(1 - x^3)^{2^\ell - 1} = (1 - x)^{2^\ell - 1} (1 + x + x^2)^{2^\ell - 1} = (1 - x)^{2^\ell - 1} \sum_{i=0}^{2(2^\ell - 1)} b_i x^i,$$

so

$$\sum_{i=0}^{2^\ell-1} x^{3i} \equiv (1-x^3)^{2^\ell-1} \equiv \frac{1-x^{2^\ell}}{1-x} \left(\sum_{i=0}^{2(2^\ell-1)} b_i x^i \right) \equiv \left(\sum_{i=0}^{2^\ell-1} x^i \right) \left(\sum_{i=0}^{2(2^\ell-1)} b_i x^i \right) \pmod{2}.$$

The part of degree $< 2^\ell$ is $\sum_{k=0}^{2^\ell-1} \beta_k x^k$. Thus $\beta_k \equiv 1 \pmod{2}$ iff $3 \mid k$. So $y_d = b_d = \beta_d - \beta_{d-1} \equiv 1 \pmod{2}$ iff $d \equiv 0, 1 \pmod{3}$. \square

Proof of Lemma 4. Write $(1+x+x^2)^{2^\ell-2} = \sum_{i=0}^{2(2^\ell-2)} c_i x^i$. Then we have $c_d = g_d$. Note that

$$(1+x+x^2) \sum_{i=0}^{2(2^\ell-2)} c_i x^i = (1+x+x^2)^{2^\ell-1} = \sum_{i=0}^{2(2^\ell-1)} b_i x^i,$$

hence $c_0 = b_0 = 1$, $c_0 + c_1 = b_1$, $c_i + c_{i-1} + c_{i-2} = b_i$, for $2 \leq i \leq 2^\ell - 2$. By Lemma 3 and by induction on i , we obtain c_i is odd if and only if $i \equiv 0, 2 \pmod{6}$. \square

Proof of Lemma 5. Since $H_j(\mathbb{N}) \subseteq \mathbb{N}$ ($1 \leq j \leq m$), we can find an integer $N \geq 1$ large enough such that for all integers $n > N$, we have $\sum_{k=i}^d \binom{k}{i} a_k^{(j)} n^{k-i} > 0$ ($0 \leq i \leq d$). Take $M > d! \cdot \sum_{k=i}^d \binom{k}{i} a_k^{(j)} n^d$, possibly depending on n . Let ℓ, b be integers such that $2^\ell > d$, $b > n$, and b is odd. Take integers z_k such that $2^{z_1} > Mb^d$, and $2^{z_k} > 2^{d z_{k-1}} Mb^d$ ($2 \leq k \leq 2^\ell$). Put $x_k = 2^{z_k}$ ($1 \leq k \leq 2^\ell$), and $A := A_n := b \sum_{k=1}^{2^\ell} x_k$.

Put $D(d, \ell) = \{\mathbf{e} = (e_k)_{1 \leq k \leq 2^\ell} \in \mathbb{N}^{2^\ell} : |\mathbf{e}| \leq d\}$. For $1 \leq j \leq m$ and $\mathbf{e} \in D(d, \ell)$, define

$$\alpha_j(\mathbf{e}) = b^{|\mathbf{e}|} C_{\mathbf{e}} \left(\sum_{k=|\mathbf{e}|}^d \binom{k}{|\mathbf{e}|} a_k^{(j)} n^{k-|\mathbf{e}|} \right) \prod_{i=1}^{2^\ell} x_i^{e_i} > 0. \tag{1}$$

Then $H_j(n) = \alpha_j(\mathbf{0})$ with $\mathbf{0} = (0, \dots, 0)$, and by multinomial expansion, we obtain further

$$H_j(n+A) = \sum_{k=0}^d a_k^{(j)} (n+A)^k = \sum_{k=0}^d a_k^{(j)} \sum_{i=0}^k \binom{k}{i} n^{k-i} b^i \left(\sum_{l=1}^{2^\ell} x_l \right)^i = \sum_{\mathbf{e} \in D(d, \ell)} \alpha_j(\mathbf{e}), \tag{2}$$

where the last summation proceeds by the colexicographic order of \mathbf{e} , which begins with $\mathbf{0}$ and ends with $(0, \dots, 0, d)$. Note that $\alpha_j(\mathbf{e}) < \mu(\alpha_j(\mathbf{f}))$ ($1 \leq j \leq m$), for all $\mathbf{f} \in D(d, \ell)$ with $\mathbf{e} < \mathbf{f}$. Indeed, if we write $\mathbf{f} = (f_k)_{1 \leq k \leq 2^\ell}$ and let k_0 be the largest index k such that $e_k < f_k$, then $e_j = f_j$ ($k_0 < j \leq 2^\ell$). Let $x_0 = 1$ if necessary, then we have

$$\begin{aligned} \mu(\alpha_j(\mathbf{f})) &\geq \prod_{k=1}^{2^\ell} x_k^{f_k} = \prod_{k=1}^{2^\ell} x_k^{f_k - e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} = x_{k_0}^{f_{k_0} - e_{k_0}} \prod_{k < k_0} x_k^{f_k - e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} \geq x_{k_0} \prod_{k < k_0} x_k^{-e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} \\ &\geq x_{k_0} x_{k_0-1}^{-\sum_{k < k_0} e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} \geq x_{k_0} x_{k_0-1}^{-d} \prod_{k=1}^{2^\ell} x_k^{e_k} > Mb^d \prod_{k=1}^{2^\ell} x_k^{e_k} \geq \alpha_j(\mathbf{e}). \end{aligned}$$

Hence the summation $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has no carry under binary expansion, since $\alpha_j(\mathbf{e}) < \mu(\alpha_j(\mathbf{f}))$. By induction on \mathbf{f} with its colexicographic order, we conclude that the binary expansion of $\sum_{\mathbf{e} < \mathbf{f}} \alpha_j(\mathbf{e})$ is a word of length $\leq v_2(\alpha_j(\mathbf{f}))$, thus the summation $\sum_{\mathbf{e} < \mathbf{f}} \alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has no carry and yields a word of equal length with that of $\alpha_j(\mathbf{f})$. So does the summation in the formula (2), hence

$$t(H_j(n+A)) = \sum_{\mathbf{e} \in D(d, \ell)} t(\alpha_j(\mathbf{e})) = t(H_j(n)) + \sum_{i=1}^d \sum_{\mathbf{e} \in D(d, \ell), |\mathbf{e}|=i} t(\alpha_j(\mathbf{e})) = t(H_j(n)),$$

since by Lemma 2, the coefficient $C_{\mathbf{e}}$ ($\mathbf{e} \neq \mathbf{0}$) appears even times in the multinomial expansion, and $t(\alpha_j(\mathbf{e})) = t(\alpha_j(\mathbf{f}))$ if $C_{\mathbf{e}} = C_{\mathbf{f}}$ and $|\mathbf{e}| = |\mathbf{f}|$. \square

Proof of Lemma 6. Let n, M be integers such that

$$da_d n + a_{d-1}^{(1)} > \frac{9}{10} \left(da_d n + a_{d-1}^{(m)} \right) > da_d, \sum_{k=i}^d \binom{k}{i} a_k^{(j)} n^{k-i} > 0,$$

and

$$M > d! \cdot \sum_{k=i}^d \binom{k}{i} a_k^{(j)} n^d \quad (0 \leq i \leq d, \text{ and } 1 \leq j \leq m).$$

Put $r = v_2(\frac{a_d d!}{\mu(a_d d!)} + 1) \geq 1$, and $q = v_2(\frac{a_d d!}{3\mu(a_d d!)} + 1)$ (if $d > 2$). Then $q = 1$ if $r > 1$ and $d > 2$. Put

$$B_j = (d-1)! \cdot \left(da_d n + a_{d-1}^{(j)} \right) / \mu(a_d d!) > 1 \quad (1 \leq j \leq m),$$

and $B_0 = \frac{9}{10} B_m$. Then

$$B_m \geq B_j \geq B_1 > B_0, \text{ for } da_d n + a_{d-1}^{(1)} > \frac{9}{10} \left(da_d n + a_{d-1}^{(m)} \right).$$

Below we shall choose appropriate $b, z_1 \in \mathbb{N}$ by distinguishing different cases.

Case 1. $d \equiv 0, 2 \pmod{6}$ and $r \equiv 1 \pmod{2}$. Then choose $b, z_1 \in \mathbb{N}$ such that $b > 16M$, $b \equiv 1 \pmod{2^{r+1}}$, and $2B_m b^{d-1} > 2^{z_1} > 2B_{m-1} b^{d-1}$, since for the integer u_0 large enough, we have

$$\begin{aligned} & \bigcup_{u \geq u_0} \left(\log_2 \left(2B_{m-1} (2^{r+1}u + 1)^{d-1} \right), \log_2 \left(2B_m (2^{r+1}u + 1)^{d-1} \right) \right) \\ & = \left(\log_2 \left(2B_{m-1} (2^{r+1}u_0 + 1)^{d-1} \right), +\infty \right), \end{aligned}$$

since

$$\lim_{u \rightarrow +\infty} \left(\log_2 \left(2B_m (2^{r+1}u + 1)^{d-1} \right) - \log_2 \left(2B_{m-1} (2^{r+1}(u+1) + 1)^{d-1} \right) \right) = \log \frac{B_m}{B_{m-1}} > 0.$$

Case 2. $d \equiv 3, 5 \pmod{6}$. Choose $b > 16M$, $b \equiv \frac{a_d d!}{\mu(a_d d!)} \pmod{4}$, and $2B_m b^{d-1} > 2^{z_1} > 2B_{m-1} b^{d-1}$.

Case 3. $d \equiv 1 \pmod{3}$, $q \equiv 1 \pmod{2}$; or $d \geq 3$, $d \equiv 2 \pmod{6}$, $r \equiv 0 \pmod{2}$. Then choose $b > 16M$, $b \equiv 1 \pmod{2^{q+r+1}}$, and $\frac{4}{3}B_m b^{d-1} > 2^{z_1} > \frac{4}{3}B_{m-1} b^{d-1}$.

Case 4. $d = 2$, $r \equiv 0 \pmod{2}$; or $d \equiv 1 \pmod{3}$, $q \equiv 0 \pmod{2}$ (thus $r = 1$). Then choose $b > 16M$, $b \equiv 1 \pmod{2^{q+r+1}}$, and $B_m b^{d-1} > 2^{z_1} > B_{m-1} b^{d-1}$.

Case 5. $d \equiv 0 \pmod{6}$, $r \equiv 0 \pmod{2}$. Take $b > 16M$, $b \equiv 1 \pmod{2^{q+r+1}}$, and $\frac{1}{2}B_m b^{d-1} > 2^{z_1} > \frac{1}{2}B_{m-1} b^{d-1}$.

Now fix $\ell \geq 1$ an integer such that $2^\ell > d$, and choose successively integers z_k ($2 \leq k \leq 2^\ell$) such that $2^{z_k} > 2^{d z_{k-1} + 2} M b^d$. Put $x_k = 2^{z_k}$ ($1 \leq k \leq 2^\ell$), $A := A_n := b \sum_{k=1}^{2^\ell} x_k$. Then

$$x_k > 4M b^d x_{k-1}^d \quad (2 \leq k \leq 2^\ell),$$

and

$$x_1 = 2^{z_1} > \frac{1}{2} B_{m-1} b^{d-1} \geq \frac{1}{2} B_1 b^{d-1} > \frac{9}{20} B_m b^{d-1} > \frac{1}{4} B_m b^{d-1} \geq \frac{1}{4} B_j b^{d-1} \quad (1 \leq j \leq m).$$

For all $\mathbf{e} = (e_j)_{1 \leq j \leq 2^\ell}$, $\mathbf{f} = (f_j)_{1 \leq j \leq 2^\ell} \in D(d, \ell)$ with $\mathbf{e} < \mathbf{f}$, we shall show below $\alpha_j(\mathbf{e}) < \alpha_j(\mathbf{f})$, and compare $\alpha_j(\mathbf{e})$ and $\mu(\alpha_j(\mathbf{f}))$ ($1 \leq j \leq m$), where $\alpha_j(\mathbf{e})$ is defined as in the formula (1).

Define $\mathbf{e}' = (e_j)_{2 \leq j \leq 2^\ell}$, $\mathbf{f}' = (f_j)_{2 \leq j \leq 2^\ell}$. For $1 \leq j \leq m$, we distinguish different cases below.

Case a. $e' < f'$. Then $\mu(\alpha_j(\mathbf{f})) > 4\alpha_j(\mathbf{e})$. Indeed if k_0 is the largest index k such that $e_k < f_k$, then $e_j = f_j$ ($k_0 < j \leq 2^\ell$), and by the construction of M, b , and x_k 's, we obtain

$$\begin{aligned} \alpha_j(\mathbf{f}) &\geq \mu(\alpha_j(\mathbf{f})) \geq \prod_{k=1}^{2^\ell} x_k^{f_k} = x_{k_0}^{f_{k_0}-e_{k_0}} \prod_{k < k_0} x_k^{f_k-e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} \geq x_{k_0} \prod_{k < k_0} x_k^{-e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} \\ &\geq x_{k_0} x_{k_0-1}^{-\sum_{k < k_0} e_k} \prod_{k=1}^{2^\ell} x_k^{e_k} \geq x_{k_0} x_{k_0-1}^{-d} \prod_{k=1}^{2^\ell} x_k^{e_k} > 4Mb^d \prod_{k=1}^{2^\ell} x_k^{e_k} \geq 4\alpha_j(\mathbf{e}). \end{aligned}$$

Case b. $e' = f'$, and $|e| < d - 1$. Then $\alpha_j(\mathbf{f}) \geq \mu(\alpha_j(\mathbf{f})) > 4\alpha_j(\mathbf{e})$. In fact, we have $B_m > 1$, and

$$\alpha_j(\mathbf{f}) \geq \mu(\alpha_j(\mathbf{f})) \geq x_1 \prod_{k=1}^{2^\ell} x_k^{e_k} > \frac{1}{4} B_m b^{d-1} \prod_{k=1}^{2^\ell} x_k^{e_k} \geq \frac{b}{4} B_m b^{|e|} \prod_{k=1}^{2^\ell} x_k^{e_k} > 4MB_m b^{|e|} \prod_{k=1}^{2^\ell} x_k^{e_k} > 4\alpha_j(\mathbf{e}).$$

Case c. $e' = f'$, and $|e| = d - 1$. Then $f_1 - e_1 = 1$, $|f| = d$, and \mathbf{f} is the successor of \mathbf{e} in $D(d, \ell)$ (i.e., there does not exist $\mathbf{g} \in D(d, \ell)$ such that $\mathbf{e} < \mathbf{g} < \mathbf{f}$). Then by definition, we have $\frac{\alpha_j(\mathbf{f})}{\alpha_j(\mathbf{e})} > \frac{bx_1}{M} > 1$, and

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} = \frac{B_j b^{d-1} \mu((e_1 + 1)!)}{x_1 e_1!} \prod_{k=2}^{2^\ell} \frac{\mu(e_k!)}{e_k!}. \tag{3}$$

From above, we deduce that if $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has a carry, then the pair (\mathbf{e}, \mathbf{f}) belongs to the Case c.

Case 1. $d \equiv 0, 2 \pmod{6}$, and $r \equiv 1 \pmod{2}$. Then

$$2B_m b^{d-1} > x_1 > 2B_{m-1} b^{d-1}, \frac{1}{2} \times \frac{10}{9} > \frac{B_m b^{d-1}}{x_1} > \frac{1}{2} \quad \text{and} \quad \frac{1}{2} > \frac{B_j b^{d-1}}{x_1} \quad (1 \leq j < m),$$

thus for all $\mathbf{e} = (e_j)_{1 \leq j \leq 2^\ell}, \mathbf{f} = (f_j)_{1 \leq j \leq 2^\ell} \in D(d, \ell)$ with $\mathbf{e} < \mathbf{f}$, we have

$$\alpha_j(\mathbf{e}) < \mu(\alpha_j(\mathbf{f})) \quad (1 \leq j < m), \quad \text{thus} \quad t(H_j(n+A)) = t(H_j(n)) \quad (1 \leq j < m),$$

just as for Lemma 5. Now $\frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} > 1$ iff $e_1 = 1$ and $e_k \in \{0, 1, 2\}$ ($2 \leq k \leq 2^\ell$), and then

$$1 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 2, \quad \frac{\alpha_m(\mathbf{f})}{\mu(\alpha_m(\mathbf{f}))} = \frac{b^d a_d d!}{\mu(a_d d!)} \equiv 2^r - 1 \pmod{2^{r+1}},$$

and we fall in the Case c. So the summation $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ takes the form

$$*** \underbrace{011 \cdots 1}_{r} \underbrace{00 \cdots 0}_s + 1 \underbrace{** \cdots *}_s,$$

and carries exactly r times. But \mathbf{f} is the successor of \mathbf{e} in $D(d, \ell)$, so for all $\mathbf{g} \in D(d, \ell) \setminus \{\mathbf{e}, \mathbf{f}\}$, either $\mathbf{g} < \mathbf{e}$ or $\mathbf{g} > \mathbf{f}$. In the first case, we have $\alpha_m(\mathbf{g}) < \mu(\alpha_m(\mathbf{e}))$ by the Cases a and b above. In the second case, we have $\alpha_m(\mathbf{e}) < \alpha_m(\mathbf{f})$ and $\mathbf{g}' > \mathbf{f}'$, thus $4\alpha_m(\mathbf{f}) < \mu(\alpha_m(\mathbf{g}))$ by the Case a, then $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f}) < 2\alpha_m(\mathbf{f}) < \mu(\alpha_m(\mathbf{g}))$. As for the proof of Lemma 5, by induction on \mathbf{e} with its colexicographic order, we obtain that the binary expansion of $\sum_{\mathbf{g} \in D(d, \ell), \mathbf{g} < \mathbf{e}} \alpha_m(\mathbf{g})$ (resp. $\sum_{\mathbf{g} \in D(d, \ell), \mathbf{g} \leq \mathbf{f}} \alpha_m(\mathbf{g})$) is a word of length $\leq v_2(\alpha_m(\mathbf{e}))$ (resp. $\leq v_2(\alpha_m(\mathbf{h}))$, for $\mathbf{h} \in D(d, \ell)$ with $\mathbf{f} < \mathbf{h}$). So the carries of $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ affect none of the other terms in the summation $\sum_{\mathbf{g} \in D(d, \ell)} \alpha_m(\mathbf{g})$. By Lemma 3, we get $t(H_m(n+A)) = t(H_m(n)) + r y_{d-2} = t(H_m(n)) + 1$, for r is odd, and $d - 2 \equiv 0, 1 \pmod{3}$. \square

Case 2. $d \equiv 3, 5 \pmod{6}$. Then for all $\mathbf{e}, \mathbf{f} \in D(d, \ell)$ with $\mathbf{e} < \mathbf{f}$, we have $\alpha_j(\mathbf{e}) < \mu(\alpha_j(\mathbf{f}))$ ($1 \leq j < m$), thus $t(H_j(n+A)) = t(H_j(n))$ ($1 \leq j < m$). But $\frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} > 1$ iff $e_1 = 1$ and $e_k \in \{0, 1, 2\}$ ($2 \leq k \leq 2^\ell$), then we are in the Case c, and

$$1 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 2, \quad \frac{\alpha_m(\mathbf{f})}{\mu(\alpha_m(\mathbf{f}))} = \frac{b^d a_d d!}{\mu(a_d d!)} \equiv 1 \pmod{4},$$

so $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has one carry. As for the Case 1, by Lemma 3 and the fact that $d-2 \equiv 0, 1 \pmod{3}$, we obtain

$$t(H_m(n+A)) = t(H_m(n)) + y_{d-2} = t(H_m(n)) + 1.$$

□

Case 3. $d \equiv 1 \pmod{3}$, $q \equiv 1 \pmod{2}$; or $d \geq 3$, $d \equiv 2 \pmod{6}$, $r \equiv 0 \pmod{2}$. Then $\frac{4}{3}B_m b^{d-1} > x_1 > \frac{4}{3}B_{m-1} b^{d-1}$, and we only consider the Case c. If $e_k \geq 3$ for some integer $k \geq 2$, then by the formula (3), we obtain

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < \frac{3}{4} \times \frac{10}{9} \times 2 \times \frac{1}{3} < 1 \quad (1 \leq j \leq m).$$

Below we assume $e_k \in \{0, 1, 2\}$ ($2 \leq k \leq 2^\ell$).

If $e_1 = 0$ or 2, then

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 1 \quad (1 \leq j \leq m).$$

If $e_1 = 1$, then for $1 \leq j \leq m$, we have $1 < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 2$, and $\frac{\alpha_j(\mathbf{f})}{\mu(\alpha_j(\mathbf{f}))} = \frac{b^d a_d d!}{\mu(a_d d!)} \equiv 2^r - 1 \pmod{2^{r+1}}$, thus the summation $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has the form

$$*** \underbrace{011 \cdots 1}_{r} \underbrace{00 \cdots 0}_{s} + 1 \underbrace{** \cdots *}_{s}.$$

hence it yields r carries (the number of such pairs is y_{d-2}).

If $e_1 = 3$, then

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 1 \quad (1 \leq j < m), 1 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 2, \frac{\alpha_m(\mathbf{f})}{\mu(\alpha_m(\mathbf{f}))} = \frac{b^d a_d d!}{3\mu(a_d d!)} \equiv 2^q - 1 \pmod{2^{q+1}}.$$

So the summation $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has the form

$$*** \underbrace{011 \cdots 1}_{q} \underbrace{00 \cdots 0}_{s} + 1 \underbrace{** \cdots *}_{s},$$

and gives q carries. Note that the number of such pairs is $y_{d-1-e_1} = y_{d-4}$, hence as for the Case 1, we obtain, by Lemma 3,

$$\begin{aligned} t(H_j(n+A)) &= t(H_j(n)) + r y_{d-2} = t(H_j(n)) \quad (1 \leq j < m), \\ t(H_m(n+A)) &= t(H_m(n)) + r y_{d-2} + q y_{d-4} = t(H_m(n)) + 1. \end{aligned}$$

□

Case 4. $d = 2$, $r \equiv 0 \pmod{2}$; or $d \equiv 1 \pmod{3}$, $q \equiv 0 \pmod{2}$ (thus $r = 1$). So $B_m b^{d-1} > x_1 > B_{m-1} b^{d-1}$. As above, we only consider the case c, and proceed similarly. If $e_k \geq 3$ for some integer $k \geq 2$, then

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < \frac{10}{9} \times 2 \times \frac{1}{3} < 1 \quad (1 \leq j \leq m).$$

Below we suppose that $e_k \in \{0, 1, 2\}$ ($2 \leq k \leq 2^\ell$). If $e_1 = 0, 2$, then

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 1 \quad (1 \leq j < m),$$

and $1 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 2$. For $e_1 = 0$, the summation $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has r carries (the number of such pairs is y_{d-1}); for $e_1 = 2$, the summation $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has q carries (the number of such pairs is y_{d-3}).

If $e_1 = 1$, then $1 < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 2$ ($1 \leq j < m$), and $2 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 3$. For $1 \leq j < m$, the summation $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has r carries (there are y_{d-2} such pairs), while the summation $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has $r-1$ carries (there are y_{d-2} such pairs).

If $e_1 = 3$, then $1 < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 2$ ($1 \leq j \leq m$), and the summation $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has q carries (there are y_{d-4} such pairs). If $e_1 \geq 4$, then

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 1 \quad (1 \leq j \leq m).$$

As for the Case 1, we obtain, by Lemma 3,

$$\begin{aligned} t(H_j(n+A)) &= t(H_j(n)) + r y_{d-2} + q y_{d-4} = t(H_j(n)) \quad (1 \leq j < m), \\ t(H_m(n+A)) &= t(H_m(n)) + r y_{d-1} + q y_{d-3} + (r-1)y_{d-2} + q y_{d-4} = t(H_m(n)) + 1. \quad \square \end{aligned}$$

Case 5. $d \equiv 0 \pmod{6}$, and $r \equiv 0 \pmod{2}$. Then $r > 1$, and $q = 1$. So $\frac{1}{2}B_m b^{d-1} > x_1 > \frac{1}{2}B_{m-1} b^{d-1}$. As above, we only consider the Case c. If $e_k > 4$ for some $k \geq 2$ or $\exists k_1, k_2 \geq 2$ with $k_1 \neq k_2$ such that $e_{k_1}, e_{k_2} \notin \{0, 1, 2\}$, then

$$\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 1 \quad (1 \leq j \leq m).$$

Now suppose $e_k = 3$ or 4 for some unique integer $k \geq 2$. If $e_1 \neq 1$, then $\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < \frac{4}{3} \times \frac{1}{3} \times 2 \times \frac{10}{9} < 1$ ($1 \leq j \leq m$). If $e_1 = 1$, then

$$1 < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 2, \quad \text{and} \quad \frac{\alpha_j(\mathbf{f})}{\mu(\alpha_j(\mathbf{f}))} = \frac{b^d a_d d!}{3\mu(a_d d!)} \equiv 2^q - 1 \pmod{2^{q+1}} \equiv 1 \pmod{4},$$

for $q = 1$. So $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has one carry. One can check that there are $(2^\ell - 1)(g_{d-5} + g_{d-6})$ such pairs: $(2^\ell - 1)g_{d-5}$ pairs for $e_k = 3$, and $(2^\ell - 1)g_{d-6}$ pairs for $e_k = 4$.

Below we assume $e_k \in \{0, 1, 2\}$ ($2 \leq k \leq 2^\ell$). If $e_1 \geq 4$, then we have $\frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 1$ ($1 \leq j \leq m$).

If $e_1 = 3$, then for $1 \leq j \leq m$, we get

$$2 < \frac{8}{3} \times \frac{9}{10} < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < \frac{8}{3} \times \frac{10}{9} < 3, \quad \frac{\alpha_j(\mathbf{f})}{\mu(\alpha_j(\mathbf{f}))} = \frac{b^d a_d d!}{3\mu(a_d d!)} \equiv 1 \pmod{4},$$

thus the summation $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ has the form

$$***01\underbrace{00 \cdots 0}_s + 10 \underbrace{** \cdots *}_s,$$

hence no carry.

If $e_1 = 0, 2$, then $1 < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 2 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 3$ ($1 \leq j < m$). If $e_1 = 0$, then $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ ($1 \leq j < m$) has r carries, while $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has $r - 1$ carries, each of them has y_{d-1} such pairs. If $e_1 = 2$, then

$$\frac{\alpha_j(\mathbf{f})}{\mu(\alpha_j(\mathbf{f}))} = \frac{b^d a_d d!}{3\mu(a_d d!)} \equiv 1 \pmod{4} \quad (1 \leq j \leq m),$$

hence $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ ($1 \leq j < m$) has one carry, and there are y_{d-3} such pairs, while $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ does not have any carry.

If $e_1 = 1$, then

$$3 < \frac{\alpha_j(\mathbf{e})}{\mu(\alpha_j(\mathbf{f}))} < 4 < \frac{\alpha_m(\mathbf{e})}{\mu(\alpha_m(\mathbf{f}))} < 5 \quad (1 \leq j < m), \quad \frac{\alpha_j(\mathbf{f})}{\mu(\alpha_j(\mathbf{f}))} \equiv 2^r - 1 \pmod{2^{r+1}} \quad (1 \leq j \leq m).$$

Thus $\alpha_j(\mathbf{e}) + \alpha_j(\mathbf{f})$ ($1 \leq j < m$) has r carries, while $\alpha_m(\mathbf{e}) + \alpha_m(\mathbf{f})$ has $r - 2$ carries, each of them has y_{d-2} such pairs.

Finally, proceeding as for the Case 1, we obtain, by Lemma 3 and Lemma 4,

$$\begin{aligned} t(H_j(n+A)) &= t(H_j(n)) + (2^\ell - 1)(g_{d-5} + g_{d-6}) + r y_{d-1} + y_{d-3} + r y_{d-2} = t(H_j(n)) \quad (1 \leq j < m), \\ t(H_m(n+A)) &= t(H_m(n)) + (2^\ell - 1)(g_{d-5} + g_{d-6}) + (r-1)y_{d-1} + (r-2)y_{d-2} = t(H_m(n)) + 1. \quad \square \end{aligned}$$

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