

VOLKER KESSLER

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## On the minimum of the unit lattice.

PAR VOLKER KESSLER

### 1. Introduction.

Computations in lattices often require a lower bound for the minimum of the lattice, both for practical purposes and for a theoretical analysis of the algorithms, e.g. [1] and [2].

In this paper we recall two results of Dobrowolski [3] and Smyth [5] in order to get such a bound for the unit lattice.

### 2. Lower bound.

Let  $K$  be a finite extension of  $\mathbb{Q}$  of degree  $n$  with maximal order  $R$ . For  $1 \leq i \leq n$  we denote by

$$K \rightarrow K^{(i)} \subset \mathbb{C}, \quad \alpha \rightarrow \alpha^{(i)}$$

the  $n$  different embeddings of  $K$  into the field  $\mathbb{C}$  of complex numbers. The first  $r_1$  of those embeddings are real, the last  $2r_2$  embeddings are non-real and numbered such that the  $(r_1 + r_2 + i)$ th embedding is the complex-conjugation of the  $(r_1 + i)$ th embedding. Then the logarithmic map is given by

$$\text{Log} : K^* \rightarrow \mathbb{R}^r, \quad \text{Log}(\alpha) := (c_1 \log |\alpha^{(1)}|, \dots, c_r \log |\alpha^{(r)}|)$$

with the unit rank  $r = r_1 + r_2 - 1$  and

$$c_i = \begin{cases} 1 & \text{for } 1 \leq i \leq r_1 \\ 2 & \text{for } r_1 + 1 \leq i \leq r + 1. \end{cases}$$

The kernel of  $\text{Log}$  consists exactly of the roots of the unity lying in  $K$ . We define the *minimum*  $\lambda(L)$  of the *unit lattice*  $L := \text{Log}(R^*)$  by

$$\lambda(L) = \min\{ \|v\| \mid v \in L \setminus \{0\} \}$$

where  $\| \cdot \|$  denotes the Euclidean norm.

**THEOREM :** *A lower bound for the minimum  $\lambda(L)$  is given by*  
 (1)

$$\lambda(L) > \mu(K) := \sqrt{\frac{2}{r+1}} \left( \frac{1}{1200} \left( \frac{\log \log n}{\log n} \right)^3 - \frac{1}{2880000} \left( \frac{\log \log n}{\log n} \right)^6 \right)$$

which is “a bit” larger than

$$\frac{1}{\sqrt{r+1}} \frac{1}{1000} \left( \frac{\log \log n}{\log n} \right)^3.$$

Thus the inverse  $1/\lambda(L)$  is of the magnitude  $O(n^{1/2+\epsilon})$  for every  $\epsilon > 0$ .

**PROOF.** Let  $\epsilon \in R^*$  be a unit of degree  $m$  over  $\mathbb{Q}$ , which is no root of unity. Without loss of generality we can assume that  $m = n$ , because if  $\|\text{Log } \epsilon\|$  is larger than  $\mu(K')$  for a subfield  $K'$  of  $K$  it is also larger than  $\mu(K)$ .

We are interested in two subsets of the conjugates  $\epsilon^{(1)}, \dots, \epsilon^{(n)}$

$$S := \{1 \leq i \leq r+1 \mid |\epsilon^{(i)}| > 1\}$$

$$T := \{1 \leq i \leq r+1 \mid |\epsilon^{(i)}| < 1\}.$$

Since  $\epsilon$  is no root of unity  $S$  is non-empty and therefore  $T$  cannot be empty because of  $N(\epsilon) = 1$ .

We call  $\epsilon$  *reciprocal* if  $\epsilon$  is conjugate to  $\epsilon^{-1}$ , i.e. its minimal polynomial  $f(X) = X^n + a_{n-1}X^{n-1} + \dots + a_0$  satisfies

$$f(X) = X^n f\left(\frac{1}{X}\right) = a_0X^n + a_1X^{n-1} + \dots + a_{n-1}X + 1.$$

If  $\epsilon$  is non-reciprocal we know from the theorem of [5] that

$$\prod_{i \in S} |\epsilon^{(i)}|^{c_i} \geq \theta$$

where  $\theta$  is the real root of  $X^3 - X - 1$ , i.e.  $\theta \approx 1.3247$ . Thus

$$(2) \quad \sum_{i \in S} c_i \log |\epsilon^{(i)}| \geq \log \theta \approx 0.281$$

But from  $N(\epsilon) = 1$  it follows

$$(3) \quad \sum_{i \in S} c_i \log |\epsilon^{(i)}| = - \sum_{i \in T} c_i \log |\epsilon^{(i)}|.$$

The value  $c_{r+1} \log |\epsilon^{(r+1)}|$  does not occur in the norm of  $\text{Log}(\epsilon)$ . But as a consequence of (3) it does not matter if  $r + 1$  lies in  $S$  or in  $T$  and so we can assume without restriction that  $r + 1 \notin S$ . Thus

$$\begin{aligned} \|\text{Log}(\epsilon)\| &\geq \sqrt{\sum_{i \in S} (c_i \log |\epsilon^{(i)}|)^2} \\ &\geq r^{-1/2} \sum_{i \in S} (c_i \log |\epsilon^{(i)}|) \geq r^{-1/2} \log \theta > \mu(K). \end{aligned}$$

(The second inequality follows from the well known norm equivalence between 1-norm and Euclidean norm.)

For reciprocal  $\epsilon$  we know by Theorem 1 of [3] :

$$(4) \quad \prod_{i \in S} |\epsilon^{(i)}|^{c_i} > 1 + \frac{1}{1200} \left( \frac{\log \log n}{\log n} \right)^3.$$

We now use the Taylor series of the logarithm ( $|y| < 1$ ) :

$$(5) \quad \log(1 + y) = y - \frac{y^2}{2} + \frac{y^3}{3} \mp \dots > y - \frac{y^2}{2}.$$

The inequality follows directly from Lagrange's representation of the residue. Applying (5) to (4) yields

$$\sum_{i \in S} c_i \log |\epsilon^{(i)}| > \frac{1}{1200} \left( \frac{\log \log n}{\log n} \right)^3 - \frac{1}{2880000} \left( \frac{\log \log n}{\log n} \right)^6.$$

Since  $\epsilon$  is reciprocal the inverses of the conjugates of  $\epsilon$  are also conjugate to  $\epsilon$ . This implies that the numbers of conjugates outside the unit circle equals the number of conjugates inside the unit circle, i.e

$$\#S = \#T \leq \frac{r+1}{2} \leq \frac{n}{2}.$$

Again by (3) we can assume that  $r + 1 \notin S$

$$\begin{aligned} \|\text{Log}(\epsilon)\| &\geq \sqrt{\sum_{i \in S} (c_i \log |\epsilon^{(i)}|)^2} \geq \sqrt{\frac{2}{r+1}} \sum_{i \in S} c_i \log |\epsilon^{(i)}| \\ &> \sqrt{\frac{2}{r+1}} \left( \frac{1}{1200} \left( \frac{\log \log n}{\log n} \right)^3 - \frac{1}{2880000} \left( \frac{\log \log n}{\log n} \right)^6 \right) = \mu(K) \end{aligned}$$

which is larger than

$$\sqrt{\frac{2}{r+1}} \left( \frac{1}{1200} - \frac{1}{2880000} \right) \left( \frac{\log \log n}{\log n} \right)^3.$$

Because of  $\sqrt{2} \left( \frac{1}{1200} - \frac{1}{2880000} \right) \approx 0.001178$  we thus proved the lower bound.

REMARK. If the conjecture of Schinzel and Zassenhaus [5] is correct the term  $\left( \frac{\log \log n}{\log n} \right)^3$  can be substituted by a constant independent of  $n$ . This bound would be provable the best one (up to constants).

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Volker Kessler  
Siemens AG  
ZFE ST SN 5  
Otto-Hahn-Ring 6  
D-8000 München 83.